



Construction of Reed-Solomon Based Quasi-Cyclic LDPC Codes Based on Protograph

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Introduction



LDPC (Low-Density Parity Check) Codes

• LDPC codes are code family with parity-check matrix of $(\# of \ 1s) \ll (\# of \ 0s)$

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\begin{bmatrix} 1 & 0 & 1 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 1 & 1 & 1 & 0 \end{bmatrix}
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- LDPC codes first proposed by Gallager in early 1960s
- LDPC code family are known to approach Shannon limits (by MacKay).
- LDPC codes have good error correcting performance with good encoding and decoding process.
- LDPC codes are selected as error correcting codes in many communication system standards.
 e.g.) IEEE 802.11 WLAN, Broadcasting(DVB-T2, ATSC3.0), 5G communication systems...

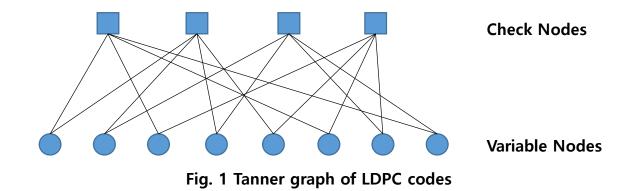


Introduction

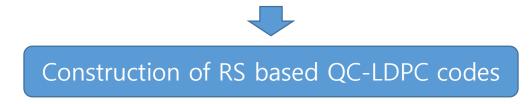


LDPC (Low-Density Parity Check) Codes

 Parity-check matrix can be represented as bipartite(Tanner) graph; Check nodes, Variable nodes, edges



- Short cycles and girths play a key role to show good BER performance
- Many researchers concerns to reduce the number of short cycles of Tanner graph





Introduction



- Protograph based LDPC codes
 - A bipartite graph with multiple edges between variable nodes and check nodes
 - We can get Tanner graph of LDPC codes using protograph with copy-and-permute procedure
 - Recently, LDPC based on protograph gains attention because protograph has simple construction and analysis.

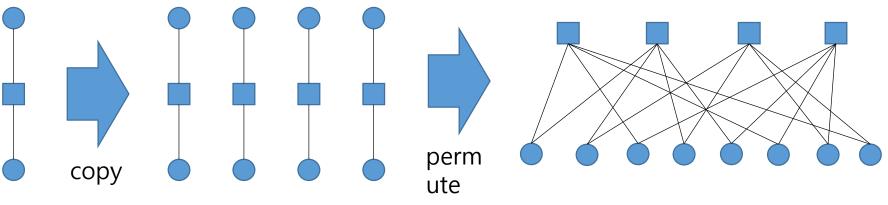
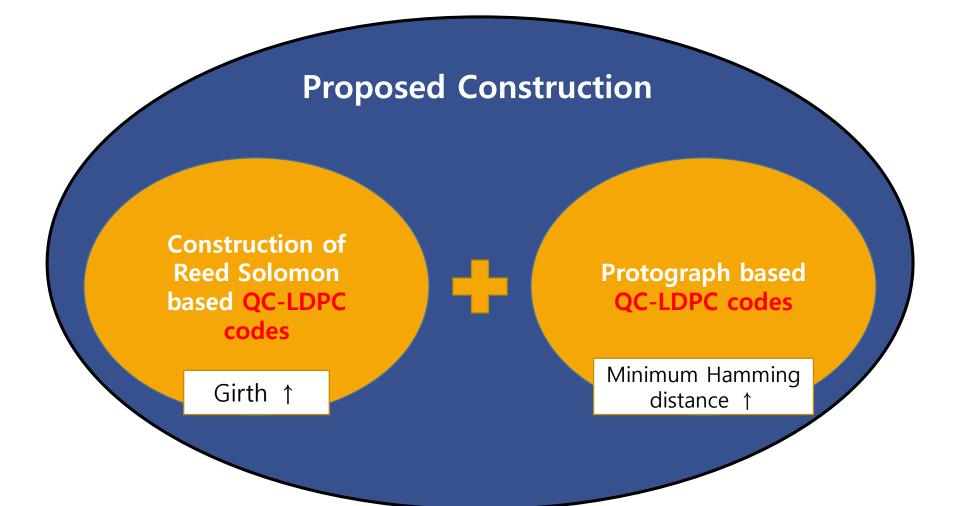


Fig. 2 Copy-and-permute procedure













Reed-Solomon based QC-LDPC codes

• Parity-check matrix of RS codes over *GF*(2^s)

$$B_{RS}(d,n) = \begin{bmatrix} 1 & \beta & \cdots & \beta^{n-1} \\ 1 & \beta^2 & \cdots & (\beta^2)^{n-1} \\ \vdots & \vdots & \ddots & \cdots \\ 1 & \beta^d & \cdots & (\beta^d)^{n-1} \end{bmatrix}$$

 β is an element of order n, where n is a factor of $2^s - 1$

• Parity-check matrix of RS based QC-LDPC codes

$$H_{RS}(d,n) = \begin{bmatrix} H_{0,0} & H_{0,1} & \cdots & H_{0,n-1} \\ H_{1,0} & H_{1,1} & \cdots & H_{1,n-1} \\ \vdots & \vdots & \ddots & \vdots \\ H_{d-1,0} & H_{d-1,1} & \cdots & H_{d-1,n-1} \end{bmatrix}$$

 $H_{i,j}$ is $r \times r$ Identity matrix cyclically shifted by the elements of $B_{RS}(i,j)$

• This RS based QC-LDPC codes has girth at least 6[1].

^[1] X. Xiao, W.E.Ryan, B.Vasic, S.Lin and K.Abdel-Ghaffar, "Reed-Solomon-Based Quasi-Cyclic LDPC Codes: Designs, Cycle Structure and Erasure Correction," Information Theory and its Application(ITA2018), Catamaran Resort, San Diego, Feb. 21-26, 2018.



Construction of RS-QC-LDPC codes



Construction of RS based QC-LDPC codes with girth 8[1]

- Make $B_{RS}(d,t)$ by choosing t columns of $B_{RS}(d,n)$
- $\Lambda_t = \{l_1, l_2, \dots, l_t\}$: Index set of selected t columns and satisfying following equations

$$\begin{split} l_{i_3} &\neq 2l_{i_2} - l_{i_1}, \ l_{i_3} \neq 3l_{i_2} - 2l_{i_1}, \ l_{i_3} \neq \frac{3l_{i_3} - l_{i_1}}{2}, \\ n &\nmid l_{i_2} + l_{i_3} - 2l_{i_1}, \ n \nmid l_{i_2} + 2l_{i_3} - 3l_{i_1}, \ n \nmid 2l_{i_2} + l_{i_3} - 3l_{i_1}, \\ n \not\restriction 2l_{i_3} - l_{i_1} - l_{i_2}, \ n \not\restriction 3l_{i_3} - 2l_{i_1} - l_{i_2}, \ n \not\restriction 3l_{i_3} - l_{i_1} - 2l_{i_2}, \end{split}$$

- Column selection is based on extending cycles of length 6 \rightarrow length 8
- RS based QC-LDPC codes with parity-check $B_{RS}(d, t)$ has girth at least 8

[1] X. Xiao, W.E.Ryan, B.Vasic, S.Lin and K.Abdel-Ghaffar, "Reed-Solomon-Based Quasi-Cyclic LDPC Codes: Designs, Cycle Structure and Erasure Correction," Information Theory and its Application(ITA2018), Catamaran Resort, San Diego, Feb. 21-26, 2018.

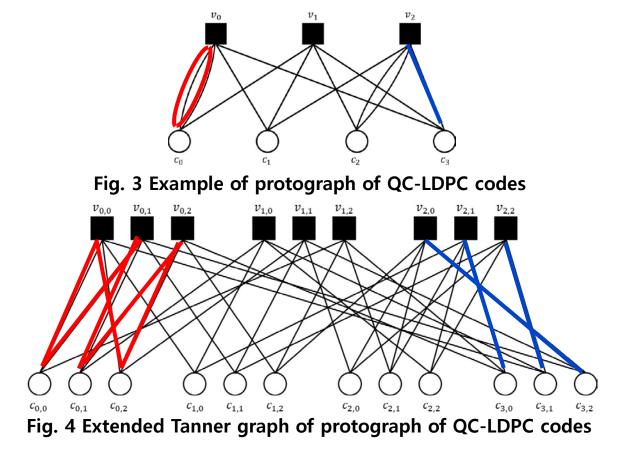


QC-LDPC codes based protograph



QC-LDPC codes based on protograph

- Protograph: allow more than 2 edges b/w check and variable nodes
- Copy small Tanner graph and lift the edges to copied graph







QC-LDPC codes based on protograph

- Calculate upper bound of minimum Hamming distance of QC-LDPC codes based on protograph[2]
- Using a function of weight matrix of protograph, they explicit the upper bound of minimum Hamming distance of protograph.
- They found out that well constructed protograph with multi-edge have better upper bound of minimum Hamming distance than protograph with single edge

$$H = \begin{bmatrix} I_0 + I_1 & I_1 & 0 & I_2 \\ I_2 & I_0 & I_1 & I_2 \\ 0 & I_1 & I_0 + I_1 & I_1 \end{bmatrix}$$

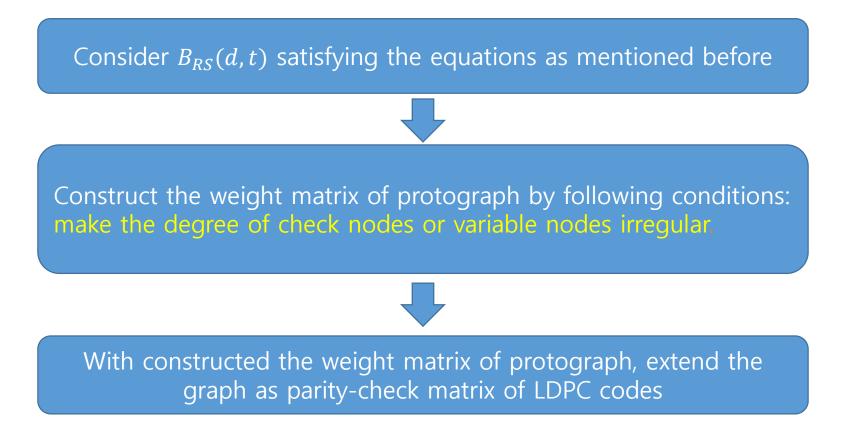
Parity matrix of Protograph of
QC-LDPC codes(Fig.4)
$$wt(H) = \begin{bmatrix} 2 & 1 & 0 & 1 \\ 1 & 0 & 1 & 1 \\ 0 & 1 & 2 & 1 \end{bmatrix}$$

Weight matrix

[2] R. Smarandache and P. O. Vontobel, "Quasi-Cyclic LDPC Codes: Influence of Proto- and Tanner-Graph Structure on Minimum Hamming Distance Upper Bounds," IEEE Trans. Info. Theory, vol. 58, no.2, Feb. 2012.



Procedures of proposed constructions





- Propose the construction of QC-LDPC codes combining RS based QC-LDPC codes with girth 8 and protograph with increasing minimum Hamming distance.
- Protograph based RS-QC-LDPC codes still holds the property that the girth is at least 8.
- Protograph based RS-QC-LDPC codes would have better errorcorrecting performance than RS based QC-LDPC codes because it has increased upper bound of minimum Hamming distance.



$$\Lambda_{8} = \{2, 5, 7, 13, 20, 32, 54, 60\} \qquad \Lambda_{12} = \\wt(H_{PRS}(4,8)) = \begin{bmatrix} 2 & 2 & 0 & 0 & 1 & 1 & 1 \\ 1 & 0 & 2 & 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & 2 & 0 & 1 \\ 1 & 1 & 1 & 1 & 0 & 0 & 2 & 2 \end{bmatrix} \qquad wt(H_{PRS}(4,8)) = Rate 1/2$$

Rate 2/3

$$\Lambda_{12} = \{0, 1, 4, 9, 11, 20, 24, 35, 41, 49, 90, 225\}$$

- Attain BER performance under sum-product decoding with maximum 50 iterations
- Use the parity-check matrix of RS codes over $GF(2^8)$
- Construct 2 weight matrix of Protograph based RS-QC-LDPC codes with rate 1/2 and 2/3.

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Simulation

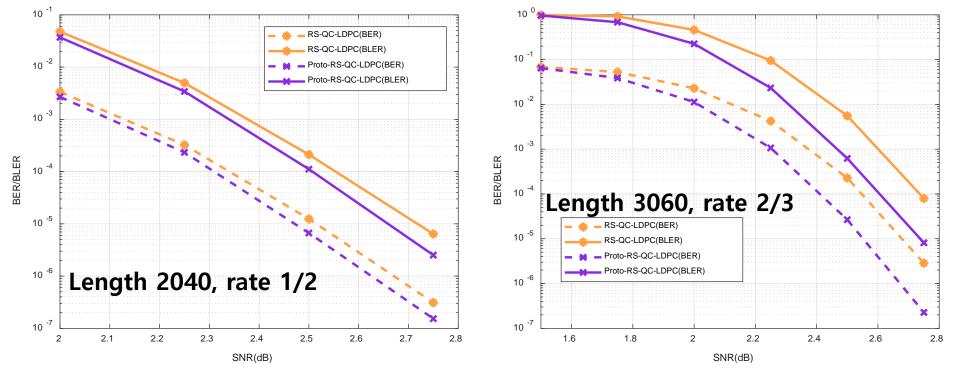


observe that BER gain is about 0.15 dB at BER 10⁻⁵ and BLER gain is about 0.12dB at BLER 10⁻⁴.

• Error correcting performance of Proto-RS-QC-LDPC have more gain when the codes have high rate.



CSDL





Conclusion



- Propose the construction of Protograph based RS-QC-LDPC codes combining two existing QC-LDPC codes construction
 - RS based QC-LDPC codes with girth 8
 - Protograph which increase the upper bounds of minimum Hamming distance
- We observe that Protograph based RS-QC-LDPC codes have better errorcorrecting performance than existing RS based QC-LDPC codes
 - We found out that the high rate codes have more coding gain.

Future work

- Can we propose protograph construction which will increase the minimum Hamming distance itself?
- Can we propose the performance analysis of Protograph based RS-QC-LDPC codes with enumerator function or EXIT chart?





Thank you for listening

Any Questions? Or comments?