



# Introduction



- LDPC codes with finite block length
    - Parameter related to performance :
      - ✓ girth (the shortest cycle)
      - ✓ connectivity of graph
    - Construction algorithms
      - ✓ ACE algorithm : connectivity of graph
      - ✓ PEG algorithm : local girth
      - ✓ **Improved PEG algorithm (IPEG) algorithm** : PEG + ACE
        - One of the best known methods for constructing LDPC codes.
  - IPEG Algorithm
    - A lot of random selection for new edges
      - ✓ About 40~50% of total edges.
- IPEG algorithm needs **more specific criterion** for selection!!

■ Stopping sets and EMD (extrinsic message degree)

EMD = 0

$A = \{v_3, v_5, v_6\}$   
Stopping set

$B = \{v_1, v_2, v_3\}$   
Not stopping set

Extrinsic check node

EMD = 2

$$\begin{bmatrix} 0 \\ 1 \\ 0 \\ 1 \end{bmatrix}_{v_3} + \begin{bmatrix} 0 \\ 1 \\ 1 \\ 0 \end{bmatrix}_{v_5} + \begin{bmatrix} 0 \\ 0 \\ 1 \\ 1 \end{bmatrix}_{v_6} = \begin{bmatrix} 0 \\ 2 \\ 2 \\ 2 \end{bmatrix}$$

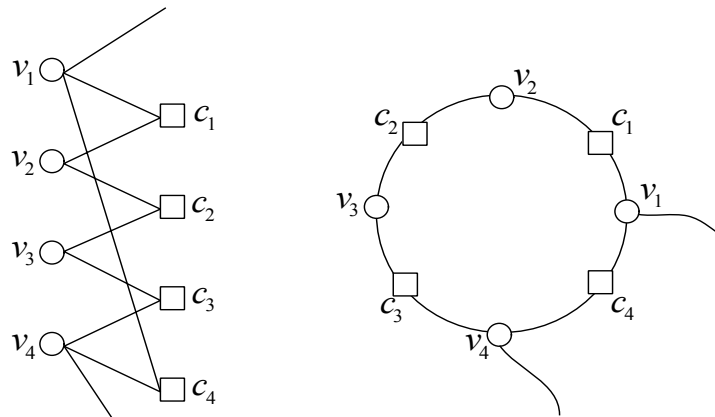
The number of 1's = 0

	$v_1$	$v_2$	$v_3$	$v_4$	$v_5$	$v_6$
$c_1$	1	1	0	1	0	0
$c_2$	0	1	1	1	1	0
$c_3$	1	0	0	0	1	1
$c_4$	0	0	1	1	0	1

$$\begin{bmatrix} 1 \\ 0 \\ 1 \\ 0 \end{bmatrix}_{v_1} + \begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \end{bmatrix}_{v_2} + \begin{bmatrix} 0 \\ 1 \\ 0 \\ 1 \end{bmatrix}_{v_3} = \begin{bmatrix} 2 \\ 2 \\ 1 \\ 1 \end{bmatrix}$$

The number of 1's = 2

- ACE (Approximate cycle EMD)
  - ACE of variable node : (degree -2), ACE of check node : 0
  - $ACE = \sum (ACE \text{ of each node})$
- EMD and ACE

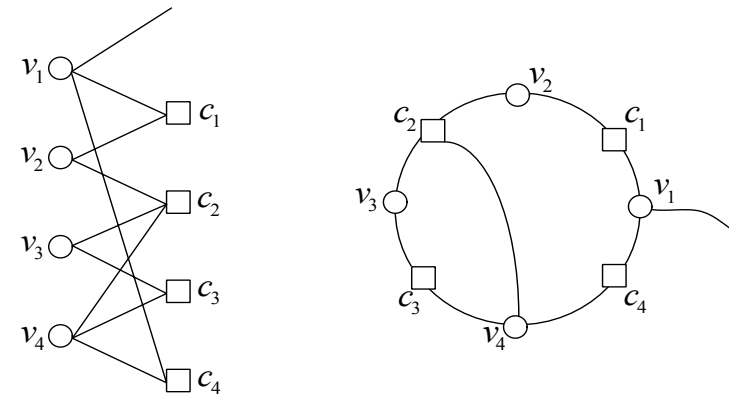


\*  $ACE = 2 = (3-2)+(2-2)+(2-2)+(3-2)$

\*  $EMD = 2$

\*  $ACE = EMD$

(a) No subcycle



\*  $ACE = 2 = (3-2)+(2-2)+(2-2)+(3-2)$

\*  $EMD = 1$

\*  $ACE \neq EMD$

(b) Subcycles



# PEG algorithm and IPEG algorithm



- Given parameter
  - the number of variable node ( $n$ ), the number of check node ( $m$ )
  - variable node degree sequence : density evolution.
- Basic idea
  - make local girth as large as possible whenever placing a new edge.

	PEG	IPEG
Construction method	Tree spreading	Tree spreading
Criterion for selection - 1st	Highest depth	Highest depth
Criterion for selection - 2nd	Lowest degree	Lowest degree
Criterion for selection - 3th	Random	ACE
Criterion for selection - 4th		Random



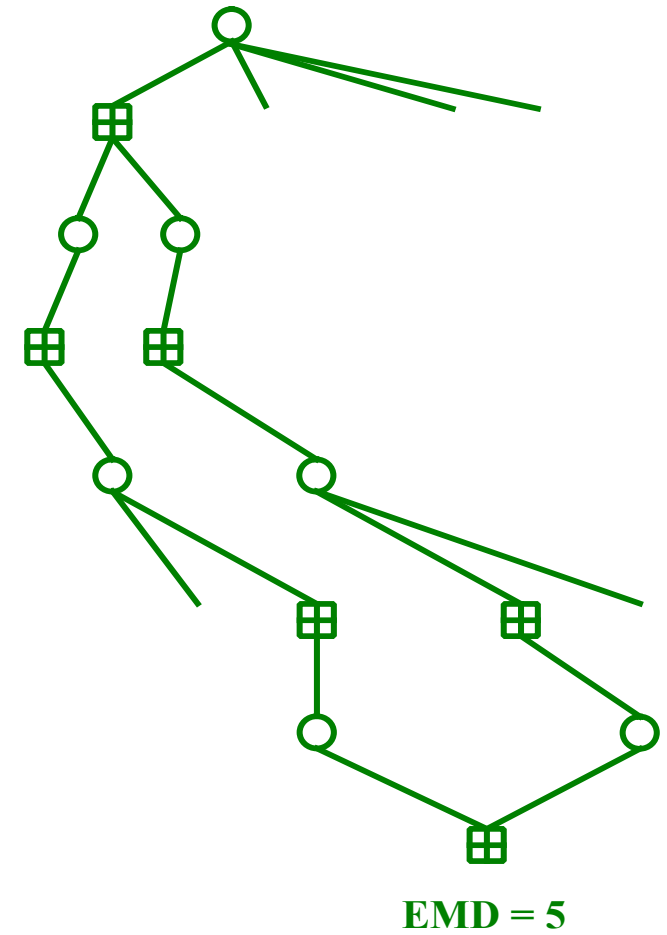
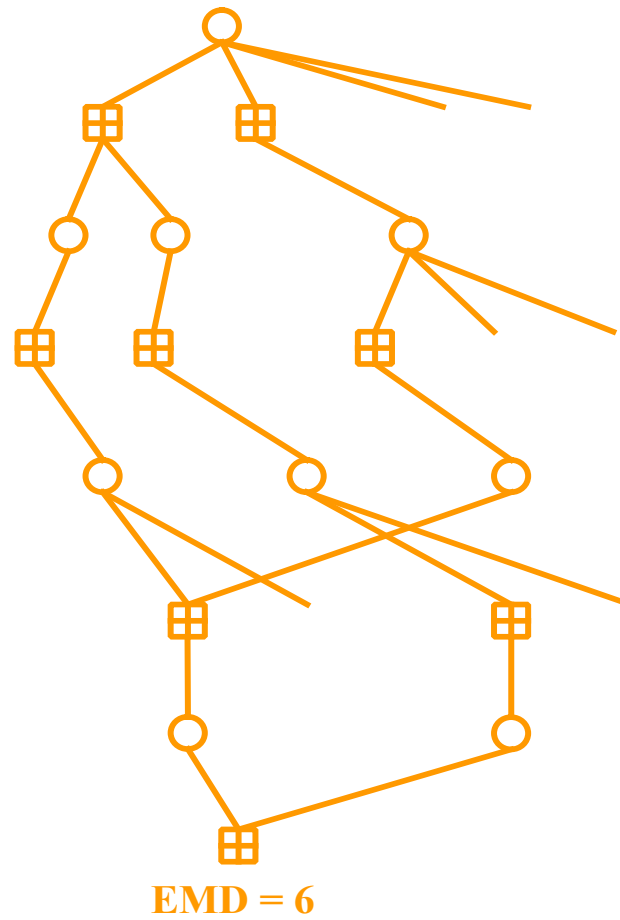
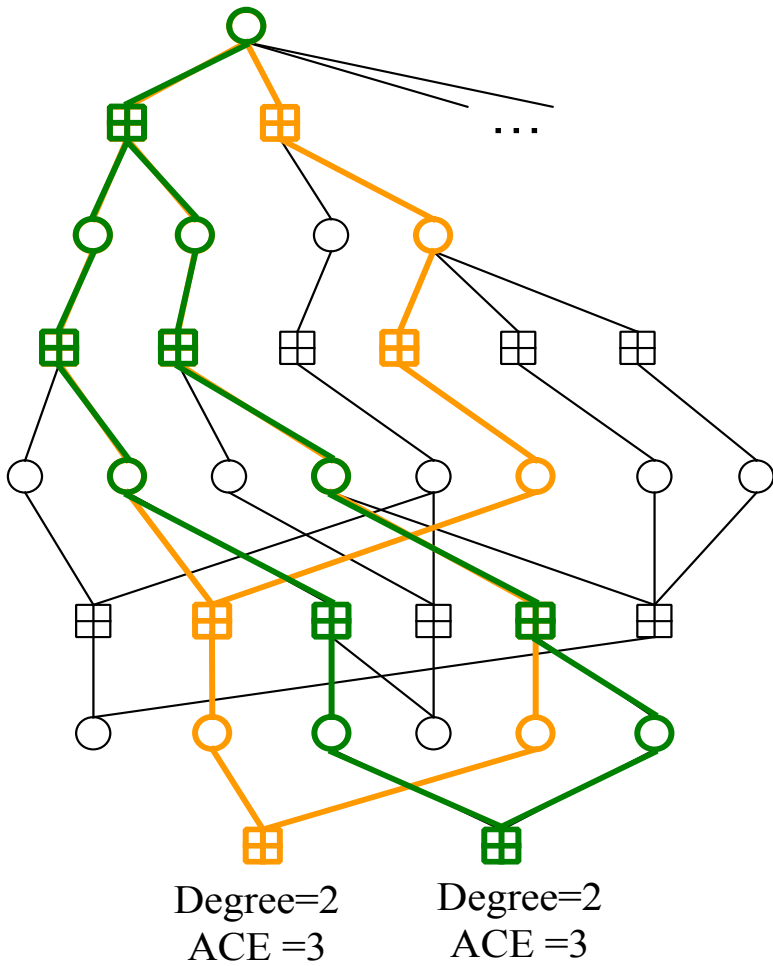
# Modified IPEG algorithm (1)



## ■ Proposed algorithm

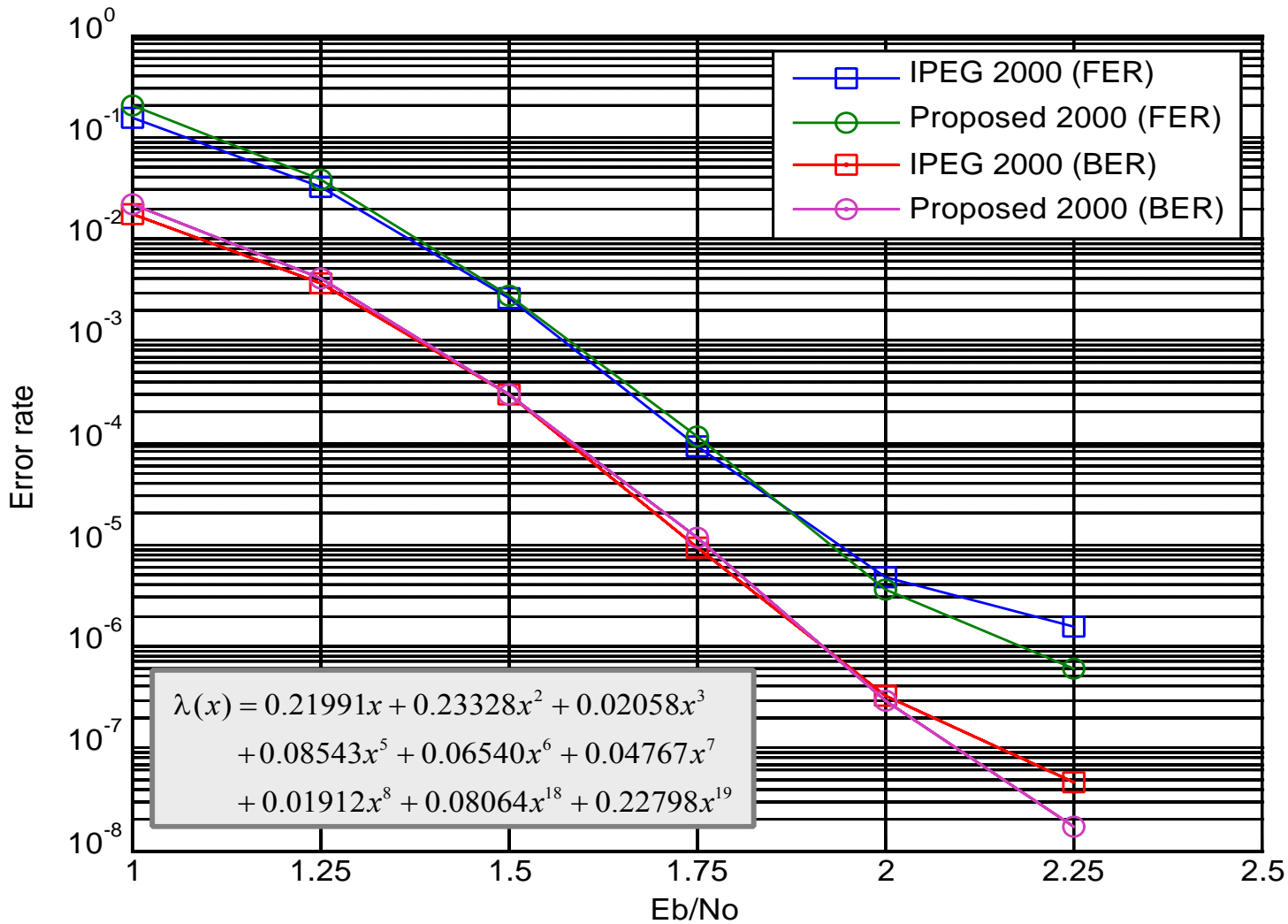
- $\Phi$  : the candidate check node classified by IPEG algorithm
- Select a check node in  $\Phi$  by calculating **the EMD of subgraph**.
  - ✓ **Extract subgraph** for each check node in  $\Phi$ 
    - By back-tracking from the check node to root (variable) node.
  - ✓ **Calculate the EMD of variable nodes in each subgraph**
    - Counting the number of 1's in the column sum over the variable nodes in subgraphs
  - ✓ Select a check node that has **the maximum EMD**
- If there are more than two check nodes that have the same EMD, select a check node that has **more variable node in subgraph**.

## ■ Example





# Simulation results (1)

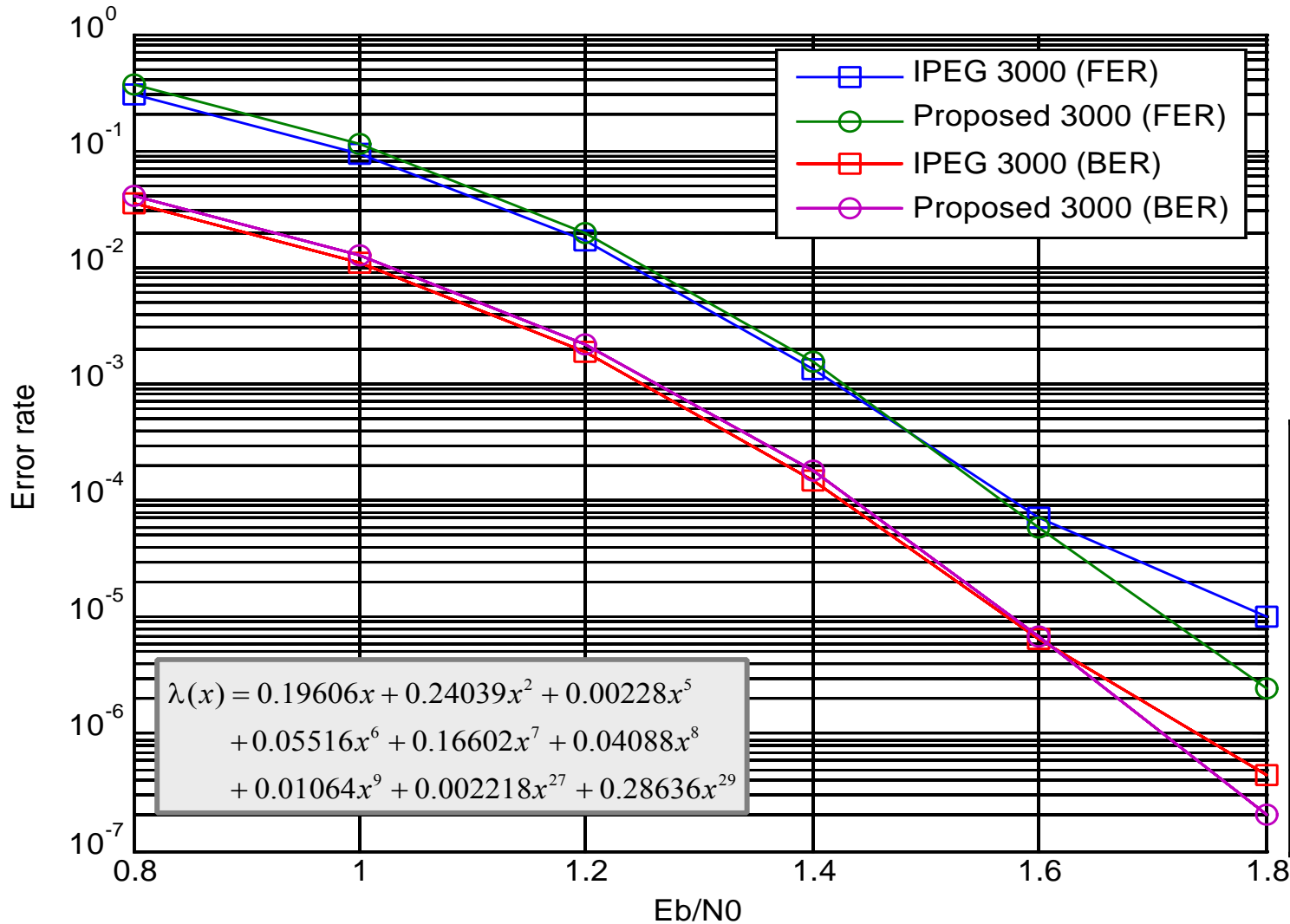


Block length = 2000  
Maximum variable node degree = 20  
Coderate 1/2

Total edges (except first)	6326
Random selection after PEG	5327 (0.84208)
Random selection after IPEG	2476 (0.43408)
Random selection after proposed	77 (0.01217)



# Simulation results (2)



Block length = 3000  
Maximum variable node degree = 30  
Coderate 1/2

Total edges (except first)	10496
Random selection after PEG	8997 (0.85718)
Random selection after IPEG	5260 (0.50114)
Random selection after proposed	148 (0.01410)