# A Research on Improvement of Error Floor of ARA codes

2006. 11. 18 Tae Ui, Kim Coding and Information Theory Lab.









Motivation

### Introduction

- Protograph and ARA codes
- Construction of parity check matrix of ARA codes
- Some theoretical bases

### Proposal to improve error floor of ARA codes

- Analysis of cycles
- Proposed algorithms
- Simulation result
- Conclusion



### Motivation



### ARA codes

- Linear time encoding
- Variable code length
- Rate compatible code
- Low threshold

### High error floor of ARA codes

• Small minimum distance of ARA results in high error floor

### Goal of the paper

• Achieve improvement of error floor without loss of performance in waterfall and with a slight burden of complexity



## Introduction



### **RA codes**

- Repetition + Accumulation
- Contrsuction based on protograph
- Decoded using BP algorithm as a subclass of LDPC codes







### Protograph and Encoder of ARA code





### Parity check matrix

- Erasure bits due to intermediate bits and puncturing
- Reducing erasrue bits for decoding convergnece







### Stopping set

- A set of variable nodes whose neighbors are connected to the set at least twice
- No smaller size of stopping set than t ensures that minimum distance  $d_{\min} \ge t$

### EMD (Extrinsic message degree)

- Number of extrinsic check nodes
- A variable node set with large EMD requires additional nodes to be a stopping set





StopspingisetEMADze 6





### **Consecutive parities cycle** $L_d$





- EMD of  $L_d$ 
  - $EMD(L_d) = Degree(b_k) 2$
  - Likely to be a small size of stopping set



# **Proposal : Union EMD**



### • Union of $L_d$ 's

- Some unions of  $L_d$ 's have interconnection
- Small size and small EMD for size
- Consider union of 2 consecutive cycle sets

### EMD of union

• Deficient EMD of union  $E_{Def}(i, j) \quad EMD(L_d^i) + EMD(L_d^j) - EMD(L_d^i \cup L_{d'}^j)$ 



Union of 2 consecutive cycle sets with size 7 and EMD 3



# **Proposal : Self return distance**



### **Self return distance** $l_s$

- Smallest number of edges from  $L_d$  to  $L_d$
- Large  $l_s$  will make more bit nodes involved when  $L_d$  and other bits node make up stopping set
- New edge does not decrease  $l_s$





# **Proposal : Edge connection**



#### Criterion for new connection

- Selection of a parity bit
  - ✓ Erasure and information connectivity
- Selection of a check node
  - ✓ Erasure and information connectivity
  - ✓ Self return distance
  - ✓ EMD
  - ✓ Maximize resulting cycle length





# **Proposal : Summary of algorithm**



- 1. Identify all the consecutive parity cycles, index from small size cycles
- **2. Evaluate initial**  $E_{Def}(i, j)$  and  $l_s^{ini}$
- 3. Add new edge
  - For consecutive cycle sets with  $E_{Def}(i, j) > 0$  until  $E_{Def}(i, j) \le 0$
  - Add new edge to an every loop except already added loop







### Interleaver construction using PEG algorithm

• Local girth distribution (K=510, Rate=0.5)

Local girth	4	6	8	10
ARA code (1)	23	247	602	147
ARA code (2)	8	111	456	444

• Consecutive parity cycle set distribution

	ARA code (1)	ARA code (2)
$L_2$	4	0
$L_3$	8	5
$L_4$	7	2
$L_5$	2	8



### **Simulation results(2)**



#### ARA code (1)



Performance of proposed scheme



19 edges are added and maximum 1 for each cycle

At FER 10<sup>-4</sup>, 0.35 dB gain

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### **Simulation results(3)**



#### ARA code (2)





- 15 edges are added and maximum 1 for each cycle
- At FER 10<sup>-4</sup>, 0.15 dB gain

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### Conclusion



### Improvement of Error floor

- Analysis of cycles with some part of the dual diagonal of ARA codes
- Supplement EMD's with the consecutive parities cycles

### Some limits

- Need enough number of short consecutive parity cycles
- Cannot provide improvement when appying to a RA codes.

### Further work

• Modification for RA codes and ARA codes with small number of short consecutive parity cycles.





# Thank you

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