



On the Merit Factor Analysis and Linear Complexity of Polyphase Sequences

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- Application of Polyphase Sequences
 - Reduction of MAI in DS-CDMA or SSMA
 - Reduction of Crest factors in OFDM and CDMA systems

- The Merit Factors and Linear Complexity of Polyphase Sequences

- All of polyphase sequences are classified as to whether the length of each sequence is square integer or not.
- **Table 1** : Classification according to Construction of Polyphase Sequences

Length	Construction	Construction
Polyphase	Frank's the proposed	Golomb's Chu

- The proposed sequence will be called Moon sequence in this thesis.

□ Construction A. $N=M^2$

Polyphase sequence $S(s_1, s_2, \dots, s_N)$

$$\Rightarrow S_{nM+k+1} = e^{i\varphi_{n,k}}, \quad 0 \leq n \leq M-1 \text{ and } 0 \leq k \leq M-1$$

1) Frank sequence : $\varphi_{n,k} = 2\pi nk/M$

2) P1 sequence : $\varphi_{n,k} = -(\pi/M)(M-2n-1)(nM+k)$

3) Px sequence : $\varphi_{n,k} = \begin{cases} (\pi/M)[(M-1)/2-k](M-2n-1), & M \text{ even} \\ (\pi/M)[(M-2)/2-k](M-2n-1), & M \text{ odd} \end{cases}$

4) The proposed sequence : $\varphi_{n,k} = \begin{cases} (\pi/M)[(M-1)/2-k](M-2n-1), & M \text{ even} \\ (\pi/M)[(M-1-k)(M-2n-1)], & M \text{ odd} \end{cases}$

□ **Example 1.** Phase representation of *Construction A* for $N=9$

1) *Frank*

0	0	0
0	$\frac{2\pi}{3}$	$\frac{4\pi}{3}$
0	$\frac{4\pi}{3}$	$\frac{2\pi}{3}$

2) P_1

0	$\frac{4\pi}{3}$	$\frac{2\pi}{3}$
0	0	0
0	$\frac{2\pi}{3}$	$\frac{4\pi}{3}$

3) P_X

$\frac{\pi}{3}$	$\frac{5\pi}{3}$	π
0	0	0
$\frac{5\pi}{3}$	$\frac{\pi}{3}$	π

4) *Moon*

$\frac{4\pi}{3}$	$\frac{2\pi}{3}$	0
0	0	0
$\frac{2\pi}{3}$	$\frac{4\pi}{3}$	0

□ **Construction B.** a positive integer N

Polyphase sequence $S(s_1, s_2, \dots, s_N)$

$$\Rightarrow s_{k+1} = e^{i\varphi_{k+1}}, \quad 0 \leq k \leq N-1$$

1) Golomb sequence :

$$\varphi_{k+1} = \frac{\pi(k+1)k}{N}$$

2) P3 sequence :

$$\varphi_{k+1} = \frac{\pi k^2}{N}$$

3) P4 sequence :

$$\varphi_{k+1} = \frac{\pi(k-N)k}{N}$$

4) Chu sequence :

$$\varphi_{k+1} = \begin{cases} \pi(k+2q)k/N, & N \text{ even} \\ \pi(k+2q+1)k/N, & N \text{ odd} \end{cases}, \text{ for an integer } q$$

□ **Example 2.** Phase representation of *Construction B* for $N=9$

1) Golomb

$$\left\{ 0 \quad \frac{2\pi}{9} \quad \frac{6\pi}{9} \quad \frac{12\pi}{9} \quad \frac{2\pi}{9} \quad \frac{12\pi}{9} \quad \frac{6\pi}{9} \quad \frac{2\pi}{9} \quad 0 \right\}$$

2)

P_3

$$\left\{ 0 \quad \frac{\pi}{9} \quad \frac{4\pi}{9} \quad \pi \quad \frac{16\pi}{9} \quad \frac{7\pi}{9} \quad 0 \quad \frac{13\pi}{9} \quad \frac{10\pi}{9} \right\}$$

3)

P_4

$$\left\{ 0 \quad \frac{10\pi}{9} \quad \frac{4\pi}{9} \quad 0 \quad \frac{16\pi}{9} \quad \frac{16\pi}{9} \quad 0 \quad \frac{4\pi}{9} \quad \frac{10\pi}{9} \right\}$$

4)

Chu

$$\left\{ 0 \quad \frac{2\pi}{9} \quad \frac{6\pi}{9} \quad \frac{12\pi}{9} \quad \frac{2\pi}{9} \quad \frac{12\pi}{9} \quad \frac{6\pi}{9} \quad \frac{2\pi}{9} \quad 0 \right\}$$

- Nonperiodic Autocorrelation Function

$$\Rightarrow \varphi_{ss}(\tau) = \sum_{n=1}^{N-\tau} s_n \cdot s_{n+\tau}^*$$

- Merit Factor F and R [3]

$$\Rightarrow F = \frac{\varphi_{ss}(0)^2}{2 \sum_{\tau=1}^{N-1} |\varphi_{ss}(\tau)|^2}$$

$$\Rightarrow R = \frac{\varphi_{ss}(0)}{\max_{1 \leq \tau \leq N-1} |\varphi_{ss}(\tau)|}$$

- F and R are important criteria for a nonperiodic autocorrelation function

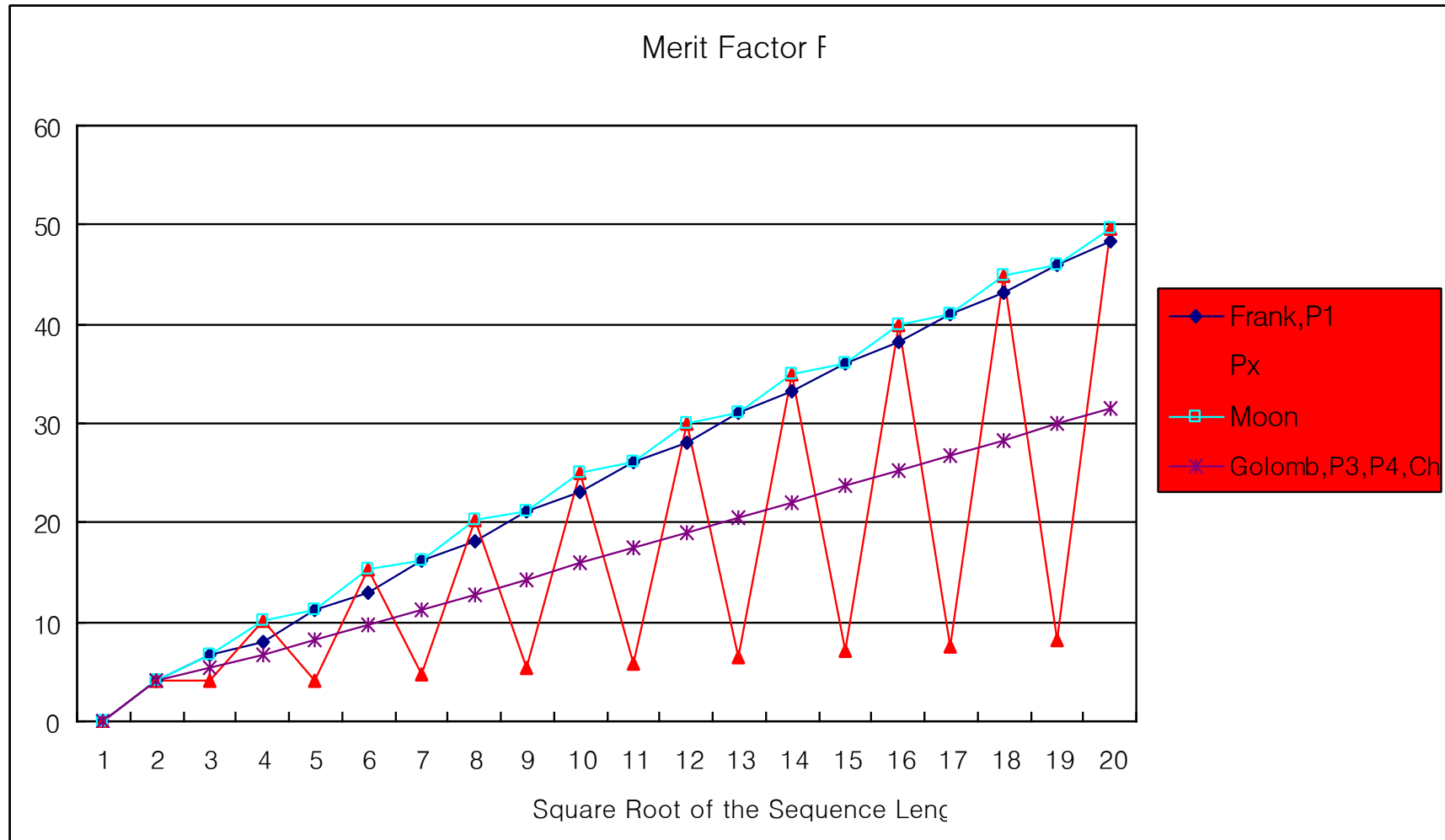


Fig 1. Merit Factor F for $\sqrt{N}$

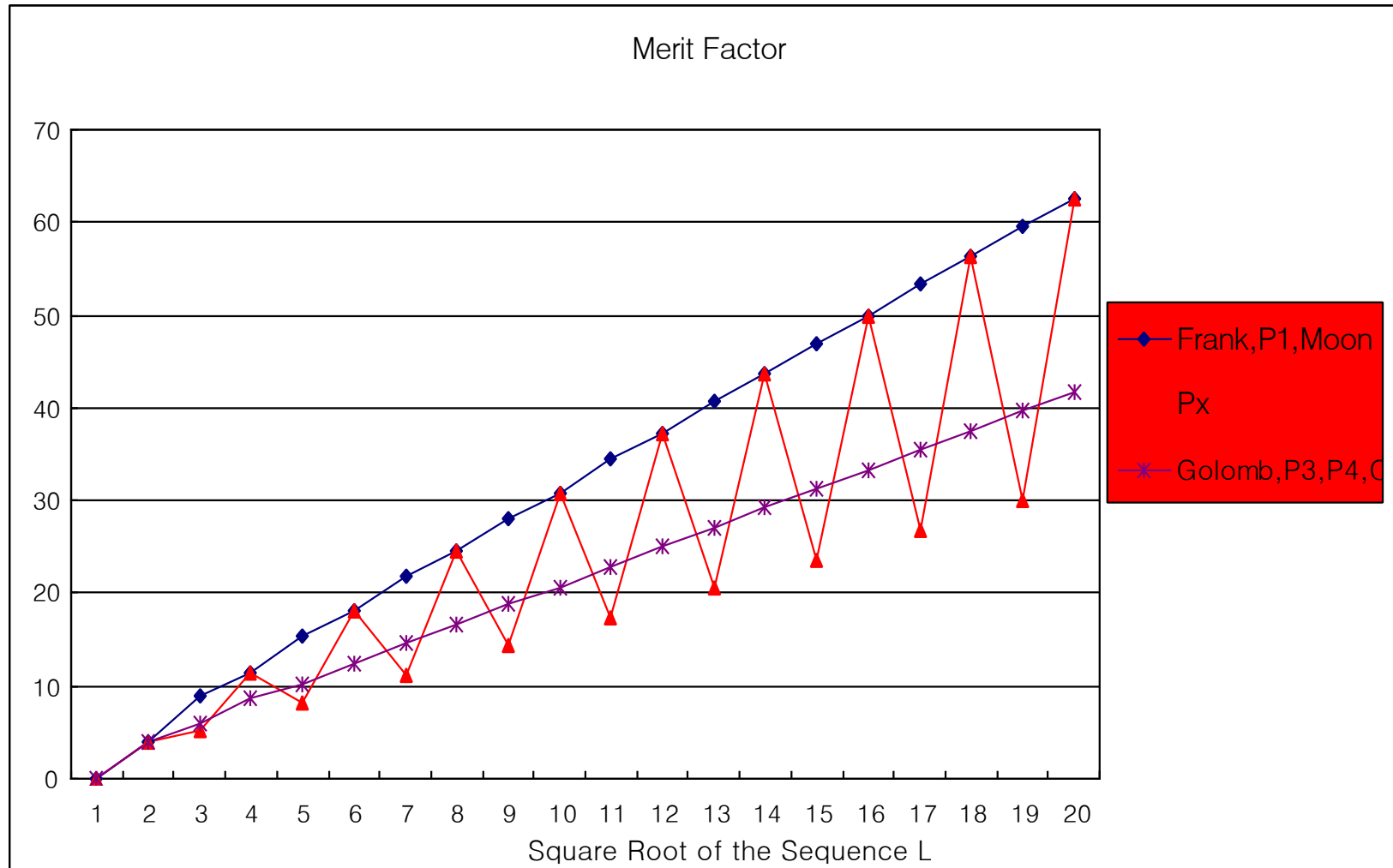


Fig 2. Merit Factor R for $\sqrt{N}$

1. LFSR Synthesis of a Polyphase Sequence over a Field

- ❑ The symbol set can't be any finite field because the set of the symbols isn't closed under addition.
- ❑ The symbol set is an complex field.
 - ⇒ an occurrence of a computation error
- ❑ A small error of the inverse operation give a huge influence on the computation recursively.

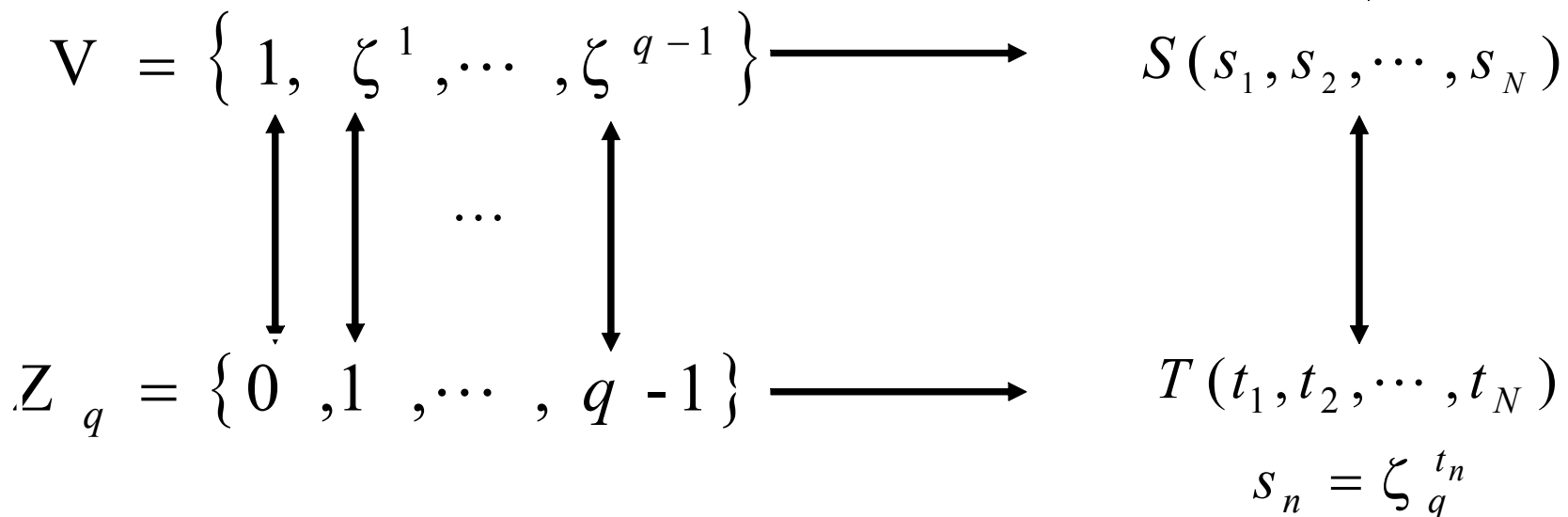
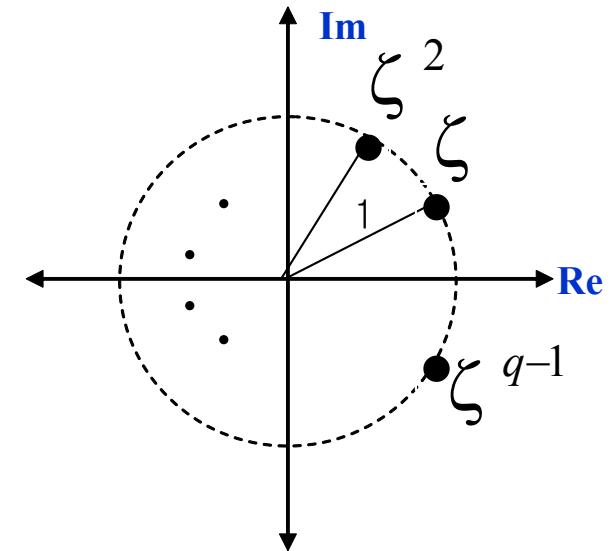
□ **Example 3.** Let the system be possible for us to calculate down to six decimal places.

- F_9 : Frank sequence S with length 9
- F_9^{new} : New sequence which is generated by LFSR

$$\begin{array}{ccccccc}
 F_9 : & \{1, 1, 1, \dots, & 1 & , & -0.5-0.866025i & , & -0.5+0.866025i\} \\
 & \downarrow \downarrow \downarrow \dots & \downarrow & & \downarrow & & \downarrow \\
 F_9^{\text{new}} : & \{1, 1, 1, \dots, & 1+0.000002i, & -0.500003-0.866026i, & -0.499998+0.866024i\}
 \end{array}$$

2. LFSR Synthesis of a Polyphase Sequence over a Ring

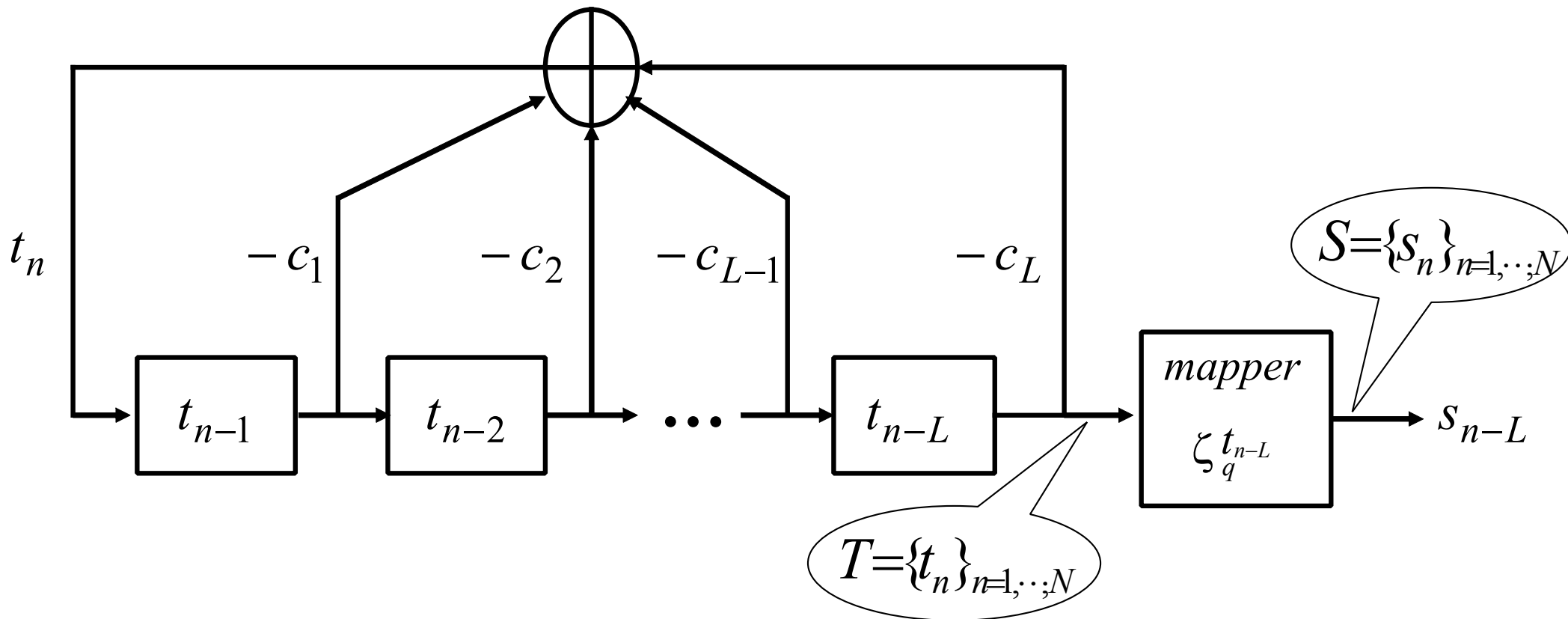
ζ_q : complex primitive q -th root of unity



□ **Table 2** : The transformed symbol set for length $N = M^2$

Original sequence	Frank	P_X	P_1	Moon
New symbol	Z_M	Z_M for N even	Z_M for N even	Z_M for N even
		Z_M for N odd	Z_M for N odd	Z_M for N odd
Original sequence	Golomb	P_3	P_4	Chu
New symbol	Z_N	Z_{2N}	Z_{2M} for N even	Z_{2N} for N even
			Z_N for N odd	Z_N for N odd

□ Polyphase sequence generator



□ The linear complexity of a polyphase sequence S can be regarded as the linear complexity of a new sequence T .

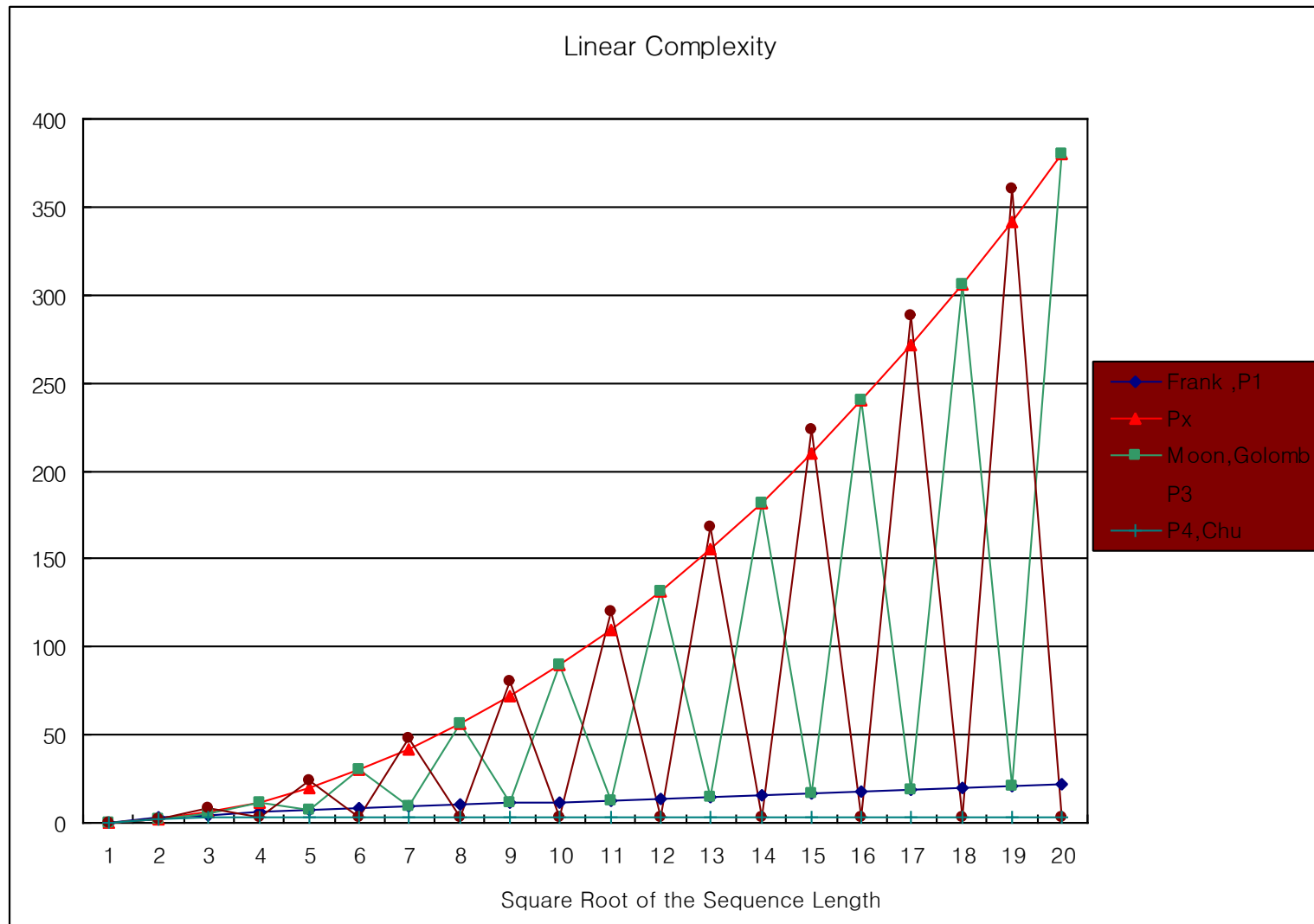


Fig 3. Linear complexities of polyphase sequences for $\sqrt{N}$

❑ Conclusion

- Merit factors analysis of each polyphase sequence
- BM-algorithm can't synthesize the shortest LFSR to generate a polyphase sequence over any field.
- Implementation of a polyphase sequence generator over a ring by the symbol set transformation
- Linear complexity analysis of each polyphase sequence by Reeds-Sloane algorithm

❑ Future Works

- Algebraic analysis of linear complexity for each polyphase sequence
- Find good polyphase sequences continuously