

Modified Concatenated LDGM codes

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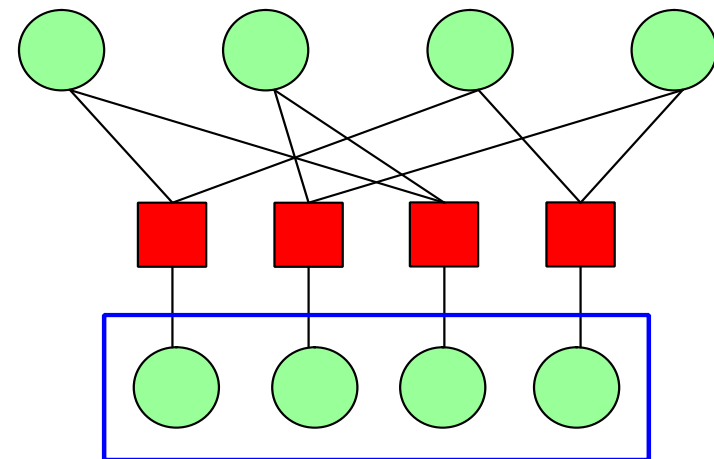


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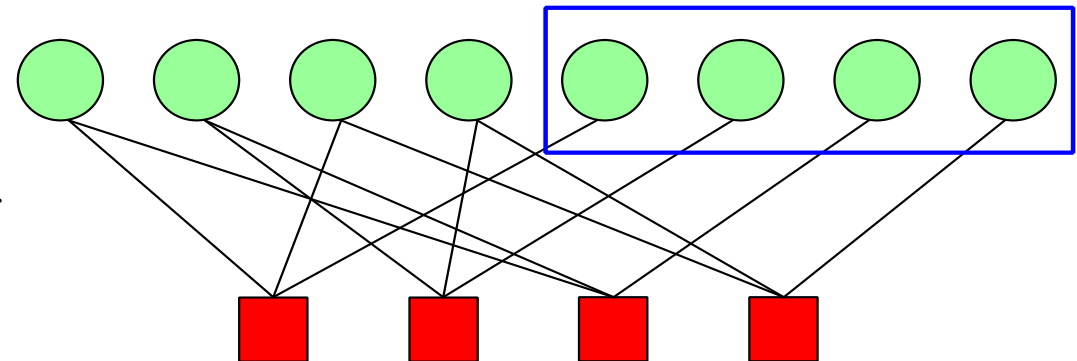
LDGM Codes

- Systematic linear block codes
- Subset of LDPC codes

$$G = \begin{bmatrix} 1 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix}$$

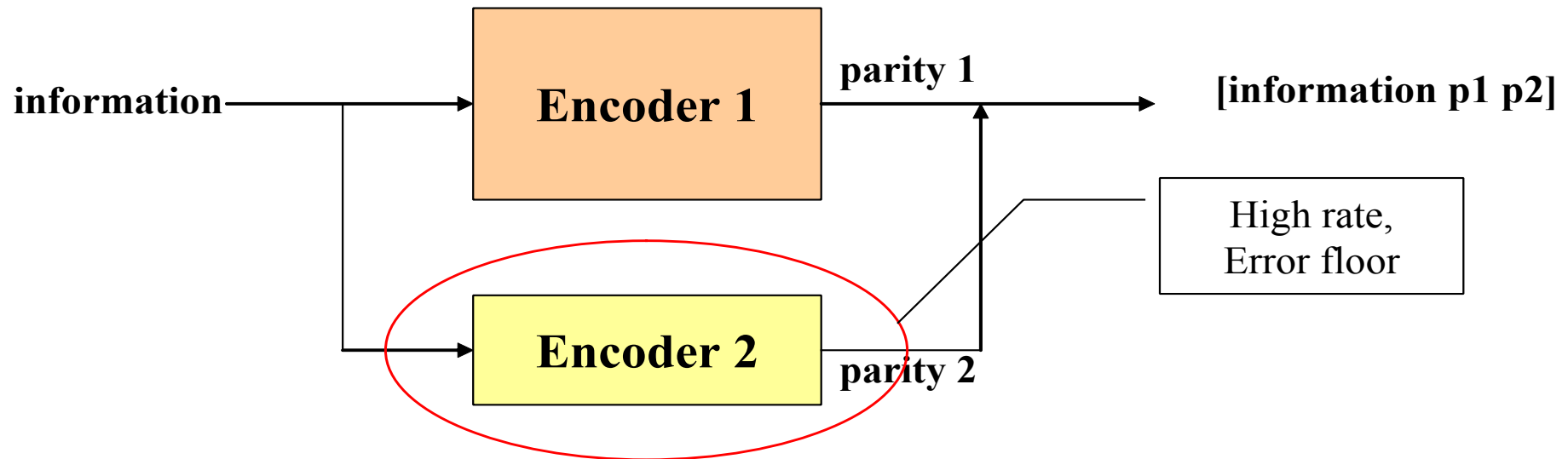


$$H = \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}$$



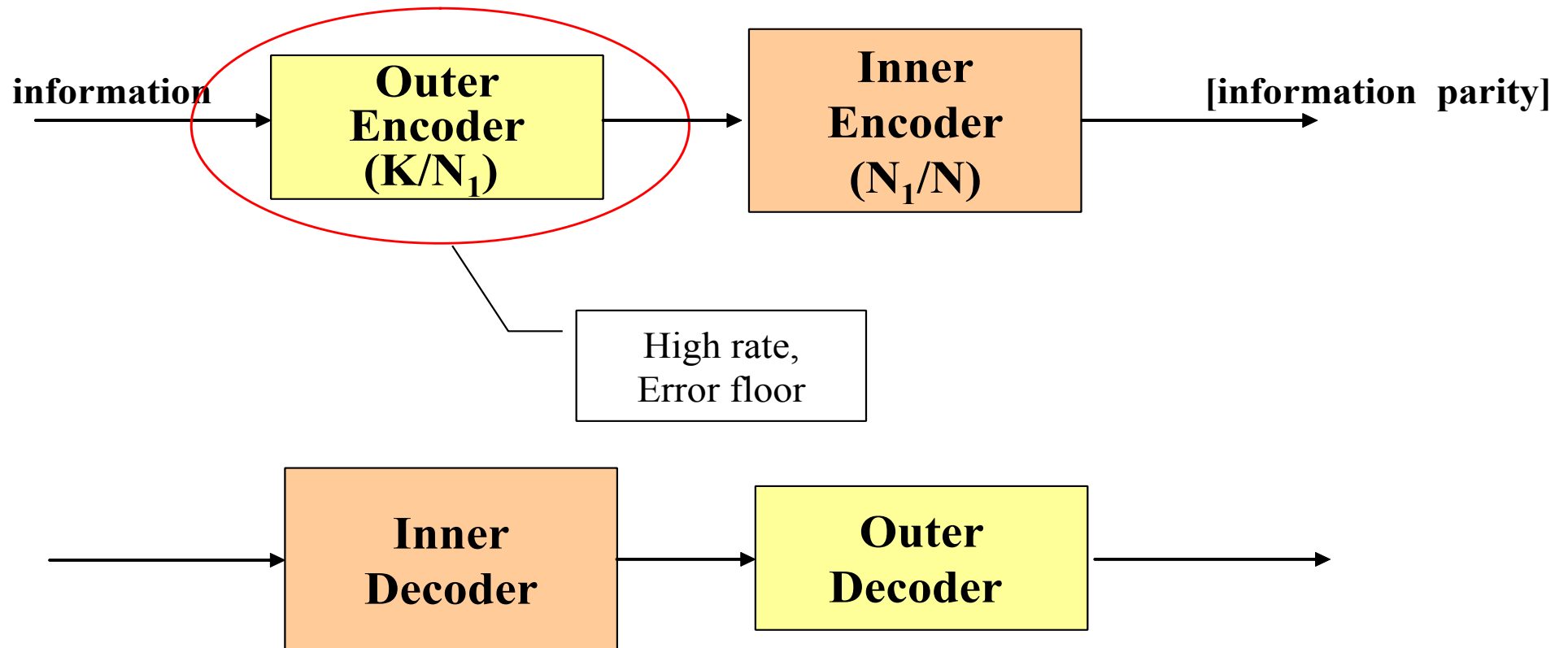
- ❑ Characteristics of LDGM Codes
 - Low encoding complexity compared to LDPC codes
 - Serious error floor
 - solution: concatenated scheme
- ❑ Serial concatenated LDGM Codes
 - Require double decoder
 - Increase the hardware complexity
- ❑ Proposed LDGM codes
 - Reduce decoder complexity about 24% without performance degradation

□ Parallel Concatenated LDGM Codes



$$G = \begin{bmatrix} \mathbf{I} & \mathbf{P}_1 & \mathbf{P}_2 \end{bmatrix} \Rightarrow H = \begin{bmatrix} \mathbf{P}_1^T & \mathbf{I} \\ \mathbf{P}_2^T & \mathbf{I} \end{bmatrix}$$

Serial Concatenated LDGM Codes



Serial Concatenated LDGM Codes

G_{OUT} : Generator matrix for outer encoder

G_{IN} : Generator matrix for inner encoder

G : Generator matrix for overall encoder

$$G = G_{OUT} \times G_{IN} = [I \ P]$$

H : Parity check matrix for decoder

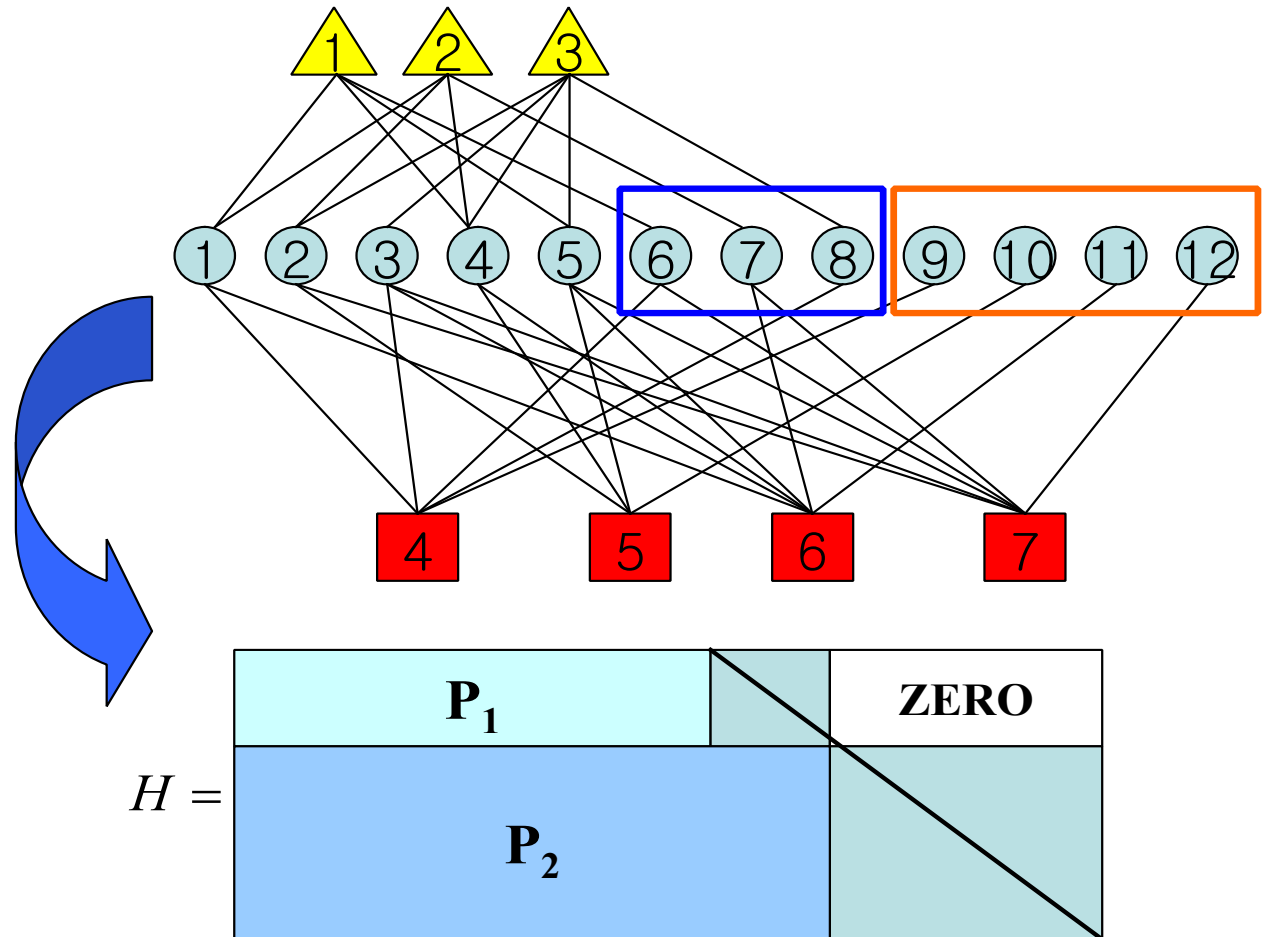
$$H = [P^T \ I]$$

- Single decoder
- Cannot be used for decoding
- Due to too many short cycles

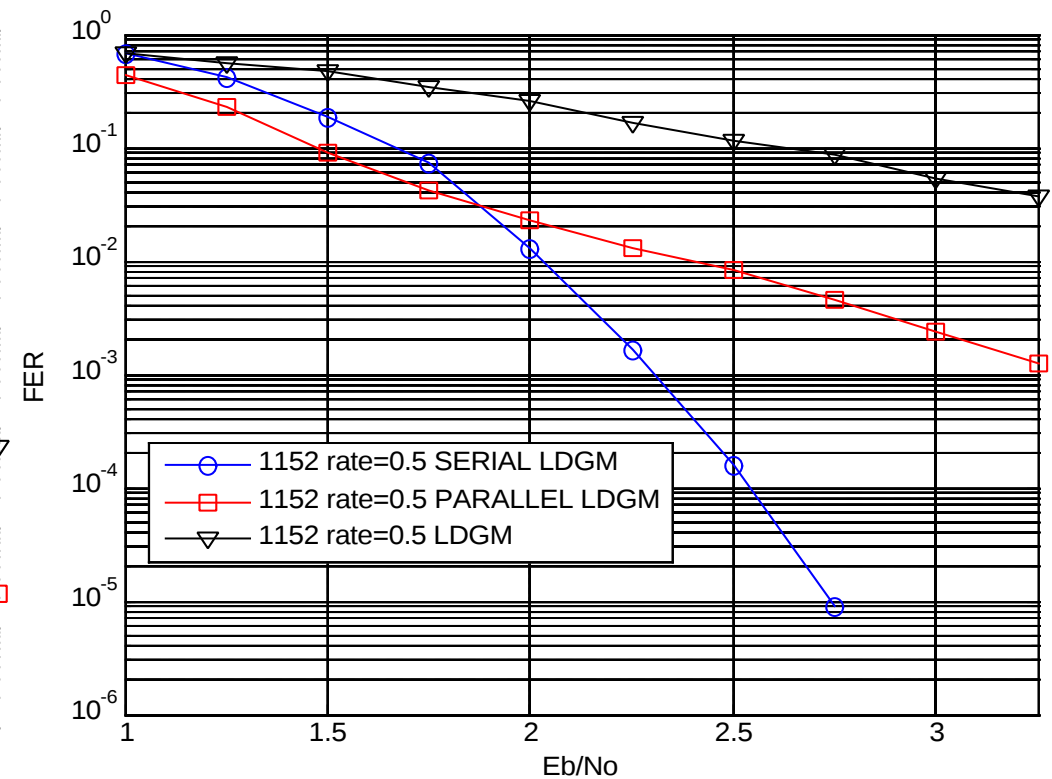
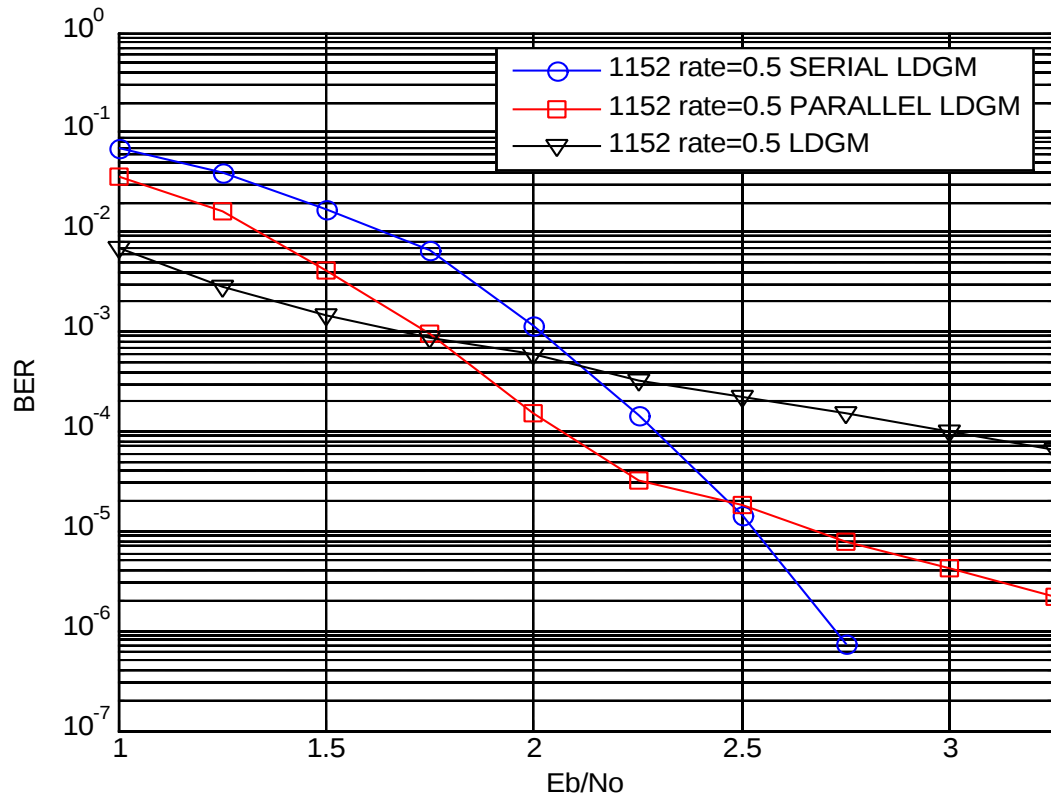
Serial Concatenated LDGM Codes

$$G_1 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 & 1 & 0 & 1 \end{bmatrix}$$

$$G_2 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$



Simulation Results



□ H matrix structure

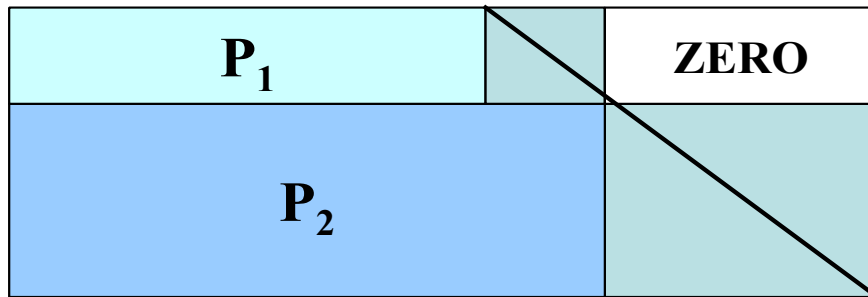


Figure 1: Parity check matrix for simply single LDGM codes.

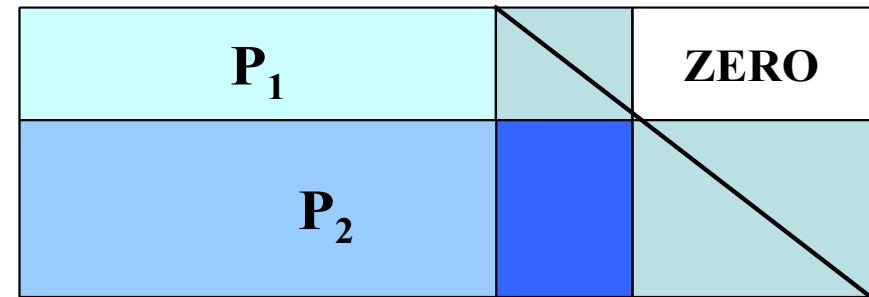
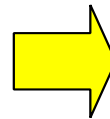


Figure 2: Parity check matrix for proposed LDGM codes.

- **To decrease density of P_1 , we lower coderate of outer encoder**
- **By decreasing the density, we reduce the number of short cycles**
- **In Figure 2, we only increase the degree of blue part to reduce the error floor**
- **By decreasing the overall degree, we reduce the complexities**

Performance of rate 0.5 proposed LDGM codes

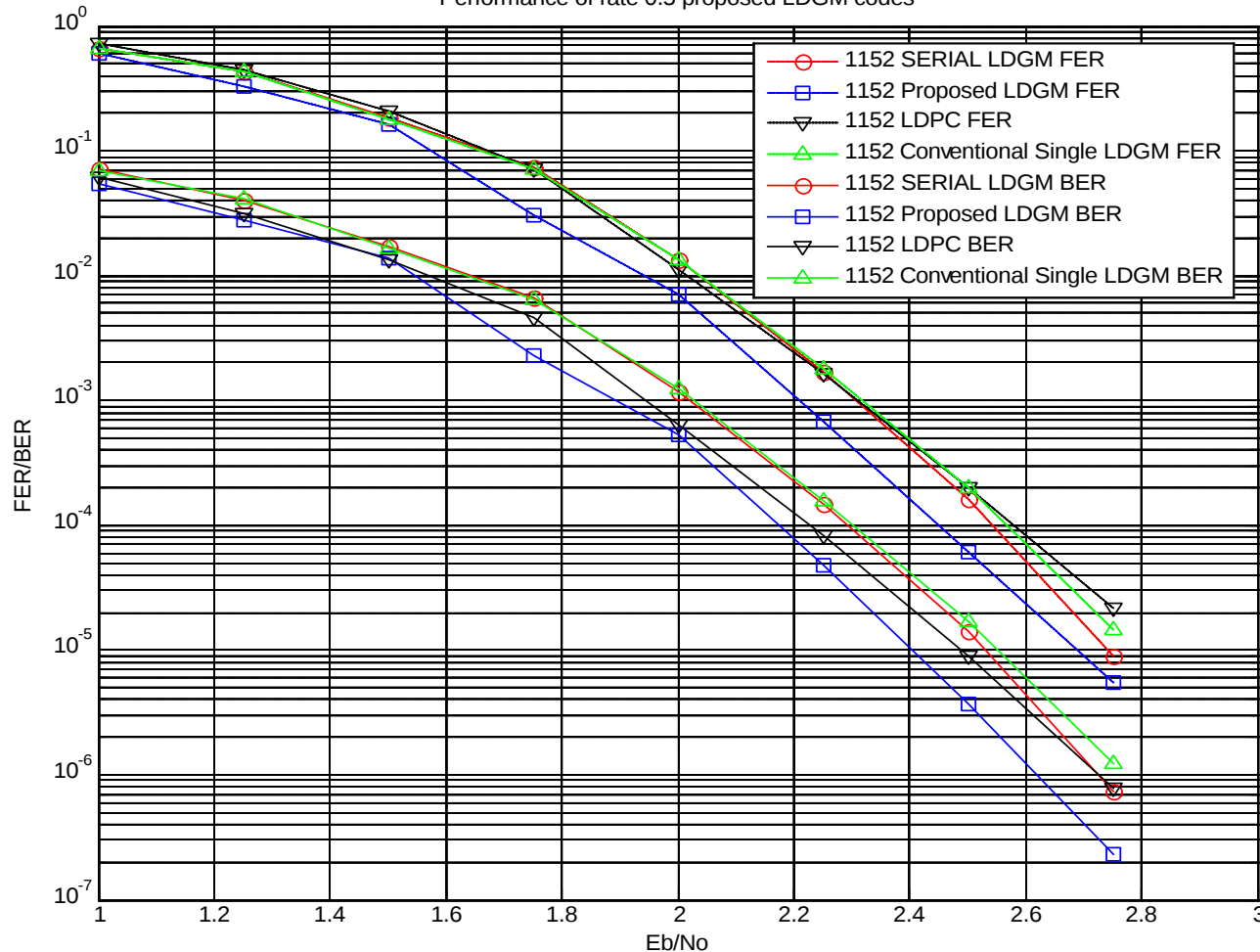


Figure 3: Performance of proposed LDGM codes.

- Block length=1152, Code rate=0.5
- Belief propagation decoding
- Maximum iteration number
 - SERIAL LDGM: 30, 50
 - Other codes: 50
- AWGN channel
- Code rate
 - SERIAL LDGM: 0.914, 0.547
 - Proposed LDGM: 0.8252, 0.6059
- PEG construction
- Number of 1's of parity check matrix

LDPC	SERIAL LDGM	Proposed LDGM
3456	6084	4642

Table 1: Number of 1's of parity check matrix.



Conclusion



- ❑ Conventional serial concatenated LDGM codes
 - Have performance similar to LDPC codes
 - Require double decoder
 - Increase decoder complexity for implementation
- ❑ Proposed LDGM codes
 - Reduce decoder complexity about 24% without performance degradation