# Trace representation and linear complexity of binary *e*-th residue sequences

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#### Zongduo Dai

State Key Laboratory of Information Security, Chinese Academy of Sciences, Beijing, China

#### **Guang Gong**

Dept. Electrical and Computer Engineering, University of Waterloo, Waterloo, ON, Canada

## Hong-Yeop Song

School of Electrical and Electronics Engineering, Yonsei University, Seoul, Korea hy.song@coding.yonsei.ac.kr

# I. Introduction

- $\diamond$  In this presentation, we would like to announce that  $\ldots$ 
  - We define binary e-th residue sequences  $s = \{s(t) | t \ge 0\}$  of period p = 1 + ef that is constant on the cosets of  $F_p^* \mod H_e$ .
  - We try to give a general description on their
    - 1. defining pairs of the form  $(g(x), \beta)$  such that  $s(t) = g(\beta^t)$  for t = 0, 1, 2, ...,
    - 2. trace representations, and
    - 3. minimal polynomials, and hence, their linear complexities.
  - For simplicity, we considered (and were able to give answers to) all e-th residue sequences for e = 2 and e = 6, and the e-th residue sequences that are characteristic sequences of e-th power residue cyclic difference sets for e = 4, 8, and 10 (as given in Baumert '71 or Storer '67 or Berndt, Evans, and Williams '98)
  - The methodology will work for any *e*-th residue sequences whether they are characteristic sequences of some cyclic difference sets or not.

• A  $(v, k, \lambda)$  cyclic difference set D is a k-subset of  $\mathbb{Z}_v \stackrel{\triangle}{=} \mathbb{Z}/v\mathbb{Z}$  such that for all non-zero  $d \in \mathbb{Z}_v$  the equation

$$x - y \equiv d \pmod{v}$$

has exactly  $\lambda$  solution pairs (x, y) with  $x, y \in D$ .

• A binary sequence  $\mathbf{s} = \{s(t) | t \ge 0\}$  (or "the characteristic sequence") of a  $(v, k, \lambda)$ -CDS of period v, defined by s(t) = 0 iff  $t \in D$ , has 2-level autocorrelation values, given as

$$\phi(\tau) = \begin{cases} v & \tau \equiv 0 \pmod{v} \\ v - 4(k - \lambda) & \tau \not\equiv 0 \pmod{v}. \end{cases}$$

• A cyclic Hadamard difference set is a (v, (v-1)/2, (v-3)/4)-cyclic difference set, and known to be equivalent to a balanced binary sequence of period v with ideal autocorrelation:  $\phi(\tau) = -1$  for all  $\tau \not\equiv 0 \pmod{v}$ .

**Conjecture 1** If a cyclic Hadamard difference set of length v exists, then v must be either

(i) a prime congruent to 3 mod 4,
(ii) a product of twin primes, or
(iii) one less than a power of 2.

- A series of computer search confirms the conjecture is true for v < 10000 except possibly for the following 13 cases: **3439**, 4355, 4623, 5775, 7395, 7743, 8227, 8463, 8591, 8835, 9135, 9215, and 9423.
  - 1. H. -Y. Song and S. W. Golomb, "On the existence of cyclic Hadamard difference sets," *IEEE Trans. Inform. Theory*, vol. 40, no. 4, pp. 1266-1268, July 1994.
  - J. -H. Kim and H. -Y. Song, "Existence of Cyclic Hadamard Difference Sets and its Relation to Binary Sequences with Ideal Autocorrelation," *Journal of Communications and Networks*, vol. 1, no.1, pp. 14-18, March 1999.
  - 3. J. -H. Kim, *On the Hadamard Sequences*, PhD Thesis, Dept Electronics Engineering, Yonsei University, Feb. 2002.

• For those three types of v, we have the following constructions:

1.  $v = p \equiv 3 \pmod{4}$  is a prime:

(a) Quadratic residue construction works for all such p.

(b) Hall's sextic residue construction works for  $p = 4x^2 + 27$ .

- 2. v = p(p+2) is a product of twin primes:
  - (a) Generalization of "Quadratic residue construction" works.

3. 
$$v = 2^t - 1$$
 for  $t = 1, 2, 3, ....$ 

- (a) m-sequence (or maximal LFSR sequence) for all such t.
- (b) GMW construction for all "composite" t.
- (c) 3-term trace sequences, 5-term trace sequences
- (d) hyperoval type (Segre Type, and Glyn Type I and Type II)

**Example 1** Binary sequences of period  $31 = 4 \cdot 1^2 + 27 = 1 + 6 \cdot 5$ . Note that 3 is a generator of  $F_{31}^*$  and we have

Cosets	Legendre	Hall's sextic
$C_* = \{0\}$		
$C_0 = \{1, 2, 4, 8, 16\}$	X	X
$C_1 = \{3, 6, 12, 24, 17\}$		X
$C_2 = \{9, 18, 5, 10, 20\}$	X	
$C_3 = \{27, 23, 15, 30, 29\}$		X
$C_4 = \{19, 7, 14, 28, 25\}$	X	
$C_5 = \{26, 21, 11, 22, 13\}$		

 The Hall's sextic residue sequence b(i) turns out to be equivalent to m-sequence of period 31 = 2<sup>5</sup> − 1.

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- Motivation of the current research
  - 1. Those of length type (i) or type (ii) are originally constructed much differently from those of length type (iii) that can naturally be described using a trace function or a sum of trace functions.
  - 2. So, what is the trace representation of those of length type (i) or (ii) ?
  - 3. What are their minimal polynomials (and hence, the linear complexity) ?
  - 4. Will it help to settle the conjecture ? Well, not much yet...

- Historical Review on "Quadratic residue sequences"
  - 1. (Turyn '64) Linear generation of quadratic residue sequences
  - 2. (Pott, '92) Abelian difference set codes
  - 3. (No, Chung, Yang, Song, '96) Trace representation of Legendre sequences of Mersenne prime period
  - 4. (Ding, Helleseth, Shan, '98) Linear complexity of Legendre sequences
  - 5. (Kim, Song, '01) Trace representation of Legendre sequences
- on "Hall's sextic residue sequences"
  - 1. (Lee, No, Chung, Yang, Kim, Song, '97) Trace representation for Mersenne Prime periods:31, 127, and 131071.
  - 2. (Kim, Song, '01) Linear complexity of Hall's sextic residue sequences
  - 3. (Kim, Gong, Song, '02) Trace representation of HSR sequences of period  $p \equiv 7 \pmod{8}$ .
  - 4. This paper completes "trace representation of HSR sequences" including the case  $p \equiv 3 \pmod{8}$
- on twin-prime sequences
  - 1. (Ding, '97) Linear complexity of generalized cyclotomic sequences of order 2
  - 2. (Kim, Song, '99) Linear complexity of binary Jacobi sequences (unpublished)
  - 3. (Dai, Gong, Song, '02) Trace representation of binary Jacobi sequences (submitted)
- The conjecture is still widely open !

# II. *e*-th residue sequences and their trace representations

- p is an odd prime, and p = ef + 1 for some e, f
- $F_p^* = F_p \setminus \{0\}$  and  $H_e = \{x^e \mid x \in F_p^*\}$
- $\alpha$  be a primitive p-th root of unity, and let  $<\alpha>^*=<\alpha>\setminus\{1\}$
- n is the order of  $2 \bmod p$ , c = (p-1)/n,  $d = \gcd(c,e)$ ,  $c_1 = c/d$ , and  $e_1 = e/d$  so that

$$ef = p - 1 = cn$$
,  $(p - 1)/d = e_1 f = c_1 n$ , and  $(e_1, c_1) = 1$ .

- $\bullet$   $LC(\mathbf{s})$  is the linear complexity of a binary sequence  $\mathbf{s}$
- $w_H(\underline{\rho})$  is the Hamming weight of a tuple  $\underline{\rho}$  over  $\overline{F}$
- $\delta(x) = 1$  (or 0) if the integer x is odd (or even), respectively.

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**Definition 1** (*e*-th residue sequences) Let  $\mathbf{s} = \{s(t) | t \ge 0\}$  be a binary sequence of period p = ef + 1. Then, we say  $\mathbf{s}$  is an *e*-th residue sequence if s(t) is constant on each of the cosets  $kH_e = \{kx \mid x \in H_e\}$  of  $H_e$  in  $F_p^*$ , that is, if  $s(t_1) = s(t_2)$  whenever  $t_1H_e = t_2H_e$ .

**Example 2 (single coset sequences)** Given any coset  $kH_e$ , let  $\mathbf{b}_{kH_e} = \{b(t)|t \ge 0\}$ , where b(t) = 1 for  $t \in kH_e$  and b(t) = 0 otherwise, then  $\mathbf{b}_{kH_e}$  is an *e*-th residue sequence.

**Example 3** Let  $\underline{1} = \{b(t) | t \ge 0\}$ , where b(t) = 1 for all t; and let  $\mathbf{b}_* = \{b(t) | t \ge 0\}$ , where b(t) = 1 if  $t = 0 \pmod{p}$  and b(t) = 0 otherwise, then these two are also *e*-th residue sequences.

We will denote the sequence  $\mathbf{b}_{kH_e}$  simply by  $\mathbf{b}_k$  with  $k \in F_p^*$ . It is clear there are only e distinct sequences in the set  $\{\mathbf{b}_k | k \in F_p^*\}$ , and they can be represented by  $\mathbf{b}_{u^i}$ , for  $0 \le i < e$ , where u is any given generator of the group  $F_p^*$ . It is clear that  $\mathbf{b}_1 = \mathbf{b}_{u^0}$  for any u, and that

$$\underline{1} = \mathbf{b}_* + \sum_{0 \le i < e} \mathbf{b}_{u^i}$$

The generating polynomial of each coset  $kH_e$  is important in expressing the trace representations of e-th residue sequences, it is defined as

$$c_{kH_e}(x) = \sum_{i \in kH_e} x^i \pmod{x^p - 1},$$

which will also be denoted simply by  $c_k(x)$  where  $k \in F_p^*$ .

**Definition 2** Given a binary sequence  $\mathbf{s} = \{s(t) | t \ge 0\}$  of period p, we say  $(g(x), \beta)$  form a defining pair of  $\mathbf{s}$  if  $s(t) = g(\beta^t)$  for t = 0, 1, 2, ..., where g(x) is a polynomial modulo  $x^p - 1$  over  $\overline{F}$  and  $\beta \in <\alpha >^*$ . We call g(x) the defining polynomial of  $\mathbf{s}$ , and  $\beta$  the corresponding defining element.

## **Theorem 1** *Let* p = ef + 1*.*

1. Let  $\mathcal{L}$  be the set of all *e*-th residue sequences of period *p*. Then  $\mathcal{L}$  is a vector space over  $F_2$  of dimension 1 + e, and  $\{\mathbf{b}_{u^i} | 0 \le i < e\} \cup \{\underline{1}\}$  is a basis of  $\mathcal{L}$  over  $F_2$ , where *u* is any given generator of  $F_p^*$ ; *i.e.*, any *e*-th residue sequence in  $\mathcal{L}$  can be expressed uniquely in the form of

$$\mathbf{s}_{\mathbf{a}^*} = a_* \underline{1} + \sum_{0 \le i < e} a_i \mathbf{b}_{u^i},$$

for some binary (1 + e)-tuple  $\mathbf{a}^* = (a_*, a_0, a_1, ..., a_i, ..., a_{e-1})$ .

2. Keep the notations in the above item, and let  $\beta \in <\alpha >^*$ . Corresponding to  $\mathbf{a}^*$  and  $\beta$ , define

$$\begin{cases} \rho_* = a_* + f \sum_{0 \le i < e} a_i, \\\\ \rho_j = \sum_{0 \le i < e} a_i c_{-u^{i+j}}(\beta) \end{cases}$$

and define

$$g(x) = \rho_* + \sum_{0 \le j < e} \rho_j c_{u^j}(x).$$

Then  $(g(x),\beta)$  is a defining pair of  $\mathbf{s}_{\mathbf{a}^*}$ .

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3. Keep the notations in the above items. Then

$$LC(\mathbf{s}_{\mathbf{a}^*}) = \delta(\rho_*) + w_H(\underline{\rho})f,$$

where

$$\underline{\rho} = (\rho_0, \rho_1, \cdots, \rho_i, \cdots, \rho_{e-1}),$$

which is given in the item 2 above.

4. Keep the notations in the above items. Let  $s_{a^*} = \{s(t) | t \ge 0\}$ . With the knowledge of the defining pair of  $s_{a^*}$  as shown in the item 2 above, its a trace representation can be obtained immediately as follows:

$$s(t) = \rho_* + \sum_{0 \le i < e} \operatorname{Tr}_1^n \left( \begin{array}{cc} \rho_i & \sum_{\substack{0 \le j < c, \\ j = i \pmod{e}}} \beta^{u^j t} \end{array} \right), \quad \forall t,$$

where  $\operatorname{Tr}_1^n(x) = \sum_{0 \le i < n} x^{2^i}$  is the trace of x from  $F_{2^n}$  to  $F_2$ .

**Proof.** The item 1 is obvious. For the item 2, we let  $r(x) = \sum_{0 \le k < p} r_k x^k \pmod{x^p - 1}$  be the defining polynomial of  $\mathbf{b}_{u^i}$  corresponding to  $\beta$ , and take the inverse Fourier transform:

$$b_{u^{i}}(t) = r(\beta^{t}) = \sum_{0 \le k < p} r_{k} \beta^{kt}, \text{ or } r_{k} = \sum_{0 \le t < p} b_{u^{i}}(t) \beta^{-kt}.$$

For k = 0,

$$r_0 = \sum_{0 \le t < p} b_{u^i}(t) = |u^i H_e| = f.$$

For  $k \in F_p^*$ , we have

$$r_{k} = \sum_{0 \le t < p} b_{u^{i}}(t)\beta^{-kt} = \sum_{t \in u^{i}H_{e}} \beta^{-kt} = \sum_{t \in -ku^{i}H_{e}} \beta^{t} = c_{-ku^{i}}(\beta).$$

Note that if  $kH_e = lH_e$  then  $-ku^iH_e = -lu^iH_e$ , and hence

$$r_k = c_{-ku^i}(\beta) = c_{-lu^i}(\beta) = r_l.$$

Therefore,  $r_k$  depends only on the coset of  $H_e$  in  $F_p^*$  to which k belongs. Denoting  $k = u^j \in F_p^*$  for j with  $0 \le j , we obtain the following useful relation:$ 

$$r_{u^j} = c_{-u^{i+j}}(\beta) = c_{-u^{i+j+em}}(\beta) = r_{u^{j+em}}, \quad \forall m.$$

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## Therefore,

$$\begin{aligned} r(x) &= f + \sum_{1 \le k < p} r_k x^k = f + \sum_{0 \le j < p-1} r_{u^j} x^{u^j} = f + \sum_{0 \le j < e} \sum_{0 \le k < f} r_{u^{j+ek}} x^{u^{j+ek}} \\ &= f + \sum_{0 \le j < e} r_{u^j} \sum_{0 \le k < f} x^{u^{j+ek}} = f + \sum_{0 \le j < e} c_{-u^{i+j}}(\beta) c_{u^j}(x) \stackrel{\triangle}{=} g_{u^i}(x). \end{aligned}$$

Clearly,  $(g_{\mathbf{1}}(x)=1,\beta)$  is a defining pair of the all-1 sequence 1. Therefore,

$$\begin{split} g(x) &= a_* + \sum_{0 \le i < e} a_i g_{u^i}(x) \\ &= a_* + \sum_{0 \le i < e} a_i \left( f + \sum_{0 \le j < e} c_{-u^{i+j}}(\beta) c_{u^j}(x) \right) \\ &= a_* + f \sum_{0 \le i < e} a_i + \sum_{0 \le j < e} \left( \sum_{0 \le i < e} a_i c_{-u^{i+j}}(\beta) \right) c_{u^j}(x) \\ &= \rho_* + \sum_{0 \le j < e} \rho_j c_{u^j}(x). \end{split}$$

The item 3 is obvious since LC is given as the number of non-zero terms in g(x). For trace representation in the item 4, we first determine the trace representation of  $\mathbf{b}_{u^i}$  as follows: Since it is a binary sequence, we have, using  $r(x) = g_{u^i}(x)$ (mod  $x^p - 1$ ),

$$\sum_{0 \le k < p} r_k^2 \beta^{2t} = r(\beta^t)^2 = b_{u^i}(t)^2 = b_{u^i}(t) = r(\beta^t) = \sum_{0 \le k < p} r_k \beta^t,$$

or

$$r_k^2 = r_{2k}$$
 or  $r_k^{2l} = r_{2^l k} \quad \forall l.$ 

Since both 2 and  $u^c$  have order n, we have  $< 2 > = < u^c >$  and

$$F_p^* = \bigcup_{0 \le i < c} u^i < u^c > = \bigcup_{0 \le i < c} u^i < 2 >.$$

Therefore, we have

$$\begin{split} b_{u^{i}}(t) &= r_{0} + \sum_{0 \leq j < p-1} r_{u^{j}} \beta^{u^{j}t} = r_{0} + \sum_{\substack{0 \leq j < c \\ 0 \leq l < n}} r_{u^{j}2^{l}} \beta^{u^{j}2^{l}t} = r_{0} + \sum_{\substack{0 \leq j < c \\ 0 \leq l < n}} \left( r_{u^{j}} \beta^{u^{j}t} \right)^{2^{l}} \\ &= r_{0} + \sum_{0 \leq j < c} \operatorname{Tr}_{1}^{n} \left( r_{u^{j}} \beta^{u^{j}t} \right) = f + \sum_{\substack{0 \leq j < c \\ 0 \leq j < c}} \operatorname{Tr}_{1}^{n} \left( c_{-u^{i+j}}(\beta) \beta^{u^{j}t} \right). \end{split}$$

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## Therefore, we have

$$\begin{split} s(t) &= a_* + \sum_{0 \le i < e} a_i b_{u^i}(t) \\ &= a_* + \sum_{0 \le i < e} a_i \left( f + \sum_{0 \le j < c} \operatorname{Tr}_1^n \left( c_{-u^{i+j}}(\beta) \beta^{u^j t} \right) \right) \\ &= a_* + f \sum_{0 \le i < e} a_i + \sum_{0 \le j < c} \operatorname{Tr}_1^n \left( \sum_{0 \le i < e} a_i c_{-u^{i+j}}(\beta) \beta^{u^j t} \right) \\ &= \rho_* + \sum_{0 \le i < e} \operatorname{Tr}_1^n \left( \rho_j \beta^{u^j t} \right) \\ &= \rho_* + \sum_{0 \le i < e} \operatorname{Tr}_1^n \left( \rho_i \sum_{\substack{0 \le j < c \\ j = i \pmod{e}}} \beta^{u^j t} \right). \end{split}$$

**Theorem 2** Let p = ef + 1, and let d be the d-parameter corresponding to the chosen (p, e). Keep the notation in Theorem 1.

1. The linear complexity of any e-th residue sequence of period p must be of the form  $\varepsilon + ke_1f$  for some  $k \in \{0, 1, 2, ..., d\}$  and  $\varepsilon \in \{0, 1\}$ . Moreover, denote by  $N_{\varepsilon+ke_1f}$  the total number of the e-th residue sequences of period p with the linear complexity being equal to  $\varepsilon + ke_1f$ . Then

$$N_{\varepsilon+ke_1f} = \begin{pmatrix} d\\ k \end{pmatrix} (2^{e_1} - 1)^k.$$

2. In the case when d = 1, we have  $N_{p-1} = N_p = 2^e - 1$ , and  $N_0 = N_1 = 1$ ; moreover, let  $\mathbf{s}_{\mathbf{a}^*}$  be the sequence as given in Theorem 1, then

$$LC(\mathbf{s_{a^*}}) = \begin{cases} p - 1 + \delta(a_* + fw_H(\mathbf{a})) & \text{if} \quad \mathbf{a} \neq (0, 0, ..., 0), \\ 1 & \text{if} \quad \mathbf{a} = (0, 0, ..., 0), \\ 0 & \text{otherwise}, \end{cases}$$

where we use the notation  $\mathbf{a} = (a_0, a_1, ..., a_{e-1})$ .

# III. *e*-tuples

Based on Theorem 1, it is enough to focus on the e-tuple of the form

$$\mathbf{c}_{u}(\beta) = (c_{u^{0}}(\beta), c_{u^{1}}(\beta), ..., c_{u^{e-1}}(\beta))$$

for trace representation and minimal polynomials, etc.

We consider the set C of the *e*-tuples  $\mathbf{c}_u(\beta)$  over all possible generators u of  $F_p^*$ and all  $\beta \in <\alpha >^*$ . That is,

$$\mathcal{C} \triangleq \{ \mathbf{c}_u(\beta) \mid \langle u \rangle = F_p^*, \ \beta \in \langle \alpha \rangle^* \}.$$

Then, it is an equivalence class under the group  $G \triangleq \langle L, D_{\lambda} | \gcd(\lambda, e) = 1, 0 < \lambda < e \} >$  where

$$L\mathbf{x} = (x_1, x_2, \cdots, x_{e-1}, x_0), \ \forall \mathbf{x} = (x_0, x_1, \dots, x_{e-1}), D_{\lambda}\mathbf{x} = (x_0, x_\lambda, x_{2\lambda}, \dots, x_{(e-1)\lambda}), \ \forall \mathbf{x} = (x_0, x_1, \dots, x_{e-1}).$$

## **Theorem 3** Let $\underline{c} = (c_0, c_1, \cdots, c_{e-1}) \in C$ , then

1.  $c_i \in F_{2^{e_1}}$  for all i.

2. <u>c</u> has  $\lambda$ -conjugate property for some integer  $\lambda$  which is coprime to  $e_1$  in the sense that

$$c_{i+dj} = c_i^{2^{\lambda j}} \quad \forall 0 \le i < e, 0 \le j < e_1.$$

Moreover, if  $\underline{c}$  has the  $\lambda$ -conjugate property, then  $D_{\nu}(\underline{c})$  has the 1-conjugate property, where  $\nu \lambda = 1 \pmod{e_1}$ .

3. Let C = (c<sub>i,j</sub>) be the square matrix of size e associated with the tuple <u>c</u>, where c<sub>i,j</sub> = c<sub>i+j</sub>, 0 ≤ i, j < e, and the index i + j are computed mod e. Then C is invertible. As a consequence, the e-tuple <u>c</u> has no smaller "period" than e. Let ε<sub>i</sub> = Tr<sub>1</sub><sup>e<sub>1</sub></sup>(c<sub>i</sub>) = ∑<sub>0≤j<e1</sub> c<sub>i</sub><sup>2j</sup>, then
(a) ε<sub>i</sub> = ∑<sub>0≤j<e1</sub> c<sub>i+dj</sub> for all i, and hence ε<sub>i+dj</sub> = ε<sub>i</sub>, for all 0 ≤ i < d, 0 ≤ j < e<sub>1</sub>.
(b) ∑<sub>0≤k<d</sub> ε<sub>k</sub> = 1,
(c) In case when d > 1, there exists at least one k in the range 0 ≤ k < d such that ε<sub>k</sub> = 0.

4. For all i = 0, 1, ..., e - 1,  $\sum_{0 \le j < e} c_j c_{j+i} = \begin{cases} f + 1 \pmod{2} & \text{if } i \equiv \frac{e\delta(f)}{2} \pmod{e} \\ f \pmod{2} & \text{otherwise,} \end{cases} \pmod{e}$ 

where the subscripts j + i are computed mod e.

5. In the case when d = 1, which is the *d*-parameter corresponding to the chosen (p, e), the *e*-tuple  $\underline{c}$  is *G*-equivalent to an *e*-tuple of the form of  $\underline{\theta} = (\theta, \theta^2, \dots, \theta^{2^{e-1}})$  for some  $\theta$ , where  $\theta$  is a root of an irreducible polynomial p(x) of degree  $e_1$  over  $F_2$ , and  $Tr_1^{e_1}(\theta) = 1$ .

## **IV.** Applications

Let p = 2f + 1 be an odd prime and u be a generator of  $F_p^*$ . Then,  $F_p = \{0\} \cup H_2 \cup uH_2$ , where  $H_2$  is the set of quadratic residues mod p and  $uH_2 = F_p^* \setminus H_2$  is the set of quadratic non-residues mod p. Let  $\mathbf{s} = \{s(t) | t \ge 0\}$  be the Legendre sequence of period p defined by the following:

$$s(t) = \begin{cases} 0 & \text{if } t \in H_2 \\ 1 & \text{otherwise.} \end{cases}$$
(1)

The item 1 of Theorem 1 implies that

$$\mathbf{s} = \underline{1} + \mathbf{b}_{u^0},$$

where <u>1</u> is the all-1 sequence. Note that  $\mathbf{a}^* = (a_*, a_0, a_1) = (1, 1, 0)$ . Therefore, from the item 3 of Theorem 1, s has a defining pair  $(g(x), \beta)$  where

$$g(x) = \rho_* + \rho_0 c_{u^0}(x) + \rho_1 c_{u^1}(x),$$

where

$$\rho_* = 1 + f, \quad \rho_j = c_{-u^j}(\beta), \quad j = 0, 1.$$

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Now, we need to determine the value of  $\mathbf{c}_u(\beta) = (c_{u^0}(\beta), c_{u^1}(\beta)) \triangleq (c_0, c_1)$ . We need the following:

**Lemma 1** Keep the notations so far. Then, the parameter d is the maximum integer that divides e and that  $x^d = 2$  has a solution in  $F_p$ .

Now, we distinguish two cases where  $2 \in H_2$  or  $2 \notin H_2$ .

Case 1  $(2 \in H_2)$ : According to the quadratic reciprocity theorem,  $2 \in H_2$ if and only if  $p \equiv 1,7 \pmod{8}$ , which are equivalent to  $f \equiv 0,3 \pmod{4}$ , respectively. This implies d = 2 from Lemma 1, and hence,  $e_1 = 2/d = 1$ . It implies that  $c_i \in F_2$  for i = 0, 1. Therefore, from the item 3 of Theorem 3,  $(\epsilon_0, \epsilon_1) = (c_0, c_1) = (1, 0)$  or (0, 1) according to the choice of u and  $\beta$ . That is,  $C = \{(1, 0), (0, 1)\}$ .

Case 2.  $(2 \in uH_2)$ : This case corresponds to  $p \equiv 3, 5 \pmod{8}$ , which are equivalent to  $f \equiv 1, 2 \pmod{4}$ , respectively. We have d = (2, c) = 1, and  $e_1 = 2/d = 2$ , and hence,  $F_2 \subset F_4 = F_{2^{e_1}} \subset F_{2^n}$ , and  $c_i \in F_4 = \{0, 1, \omega, \omega^2\}$ for i = 0, 1, where  $\omega$  is a primitive 3-rd root of unity. From Theorem 3, the fact that d = 1 implies  $\epsilon_0 = 1 = c_0 + c_1$ . Therefore,  $c_i \in F_4 \setminus F_2$  for i = 0, 1, and we have  $\mathcal{C} = \{(\omega^2, \omega), (\omega, \omega^2)\}$ . In conclusion, we may choose  $\beta \in <\alpha>^*$  such that for any given generator u of  $F_p^*$ , we have

$$(c_{u^0}(\beta), c_u(\beta)) = \begin{cases} (1,0) & \text{if } p = 1 \pmod{8} \\ (0,1) & \text{if } p = 7 \pmod{8} \\ (w^2, w) & \text{if } p = 3 \pmod{8} \\ (w, w^2) & \text{if } p = 5 \pmod{8}, \end{cases}$$

where  $\omega \in F_4$  is a primitive 3-rd root of unity. With  $\beta$  and  $\omega$  chosen as in the above,  $(g(x), \beta)$  is a defining pair of s, where

$$g(x) = \frac{p+1}{2} + \begin{cases} c_{u^0}(x) & \text{if } p = \pm 1 \pmod{8} \\ wc_{u^0}(x) + w^2c_{u^1}(x) & \text{if } p = \pm 3 \pmod{8}. \end{cases}$$

The linear complexity of s is given as

$$LC(\mathbf{s}) = \delta(\frac{p+1}{2}) + \begin{cases} \frac{p-1}{2} & \text{if } p = \pm 1 \pmod{8} \\ p-1 & \text{if } p = \pm 3 \pmod{8}. \end{cases}$$

**Theorem 4** Let p = ef + 1 be a prime with e = 6 and f odd. Let d be the d-parameter corresponding to the chosen (p, 6). Then

1. (Sextic residue sequences in general) There exist a generator u of  $F_p^*$  and  $\beta \in <\alpha >^*$  such that

$$\mathbf{c}_{u}(\beta) = \begin{cases} (0, 1, 1, 0, 1, 0) & \text{if } d = 6, \\ (1, 0, w, 1, 0, w^{2}) & \text{if } d = 3, \\ (\gamma, \gamma^{3}, \gamma^{2}, \gamma^{6}, \gamma^{4}, \gamma^{5}) & \text{if } d = 2, \\ (\theta, \theta^{2}, \theta^{4}, \theta^{8}, \theta^{16}, \theta^{32}) & \text{if } d = 1, \end{cases}$$

where

- w is a root of  $x^2 + x + 1$ ,
- $\gamma$  is a root of  $x^3 + x + 1$ , and
- $\theta = \rho$  or  $\theta = \rho + 1$  where  $\rho$  is a root of  $x^6 + x^5 + 1$  (and hence,  $\rho + 1$  is a root of  $x^6 + x^5 + x^2 + x + 1$ ).
- (Hall's sextic residue sequences) In the case when p = 6f + 1 = 4z<sup>2</sup> + 27 for some integer z, let s be the Hall's sextic residue sequence of period p which is defined as the characteristic sequence of the Hall's sextic residue different set D = H<sub>6</sub> ∪ u<sup>3</sup>H<sub>6</sub> ∪ u<sup>i</sup>H<sub>6</sub>, where u<sup>i</sup>H<sub>6</sub> is the coset containing 3. Then

(a) There exists a generator u of  $F_p^*$  and  $\beta \in <\alpha>^*$  such that

$$\mathbf{c}_u(\beta) = \begin{cases} (0, 1, 1, 0, 1, 0) & \text{if } p = 7 \pmod{8} \\ (1, 0, w, 1, 0, w^2) & \text{if } p = 3 \pmod{8} \end{cases}$$

(b) With the choice of u and  $\beta$  as in the above item,  $(g(x), \beta)$  is a defining pair of s, where

$$g(x) = \begin{cases} c_{u^0}(x) & \text{if } p = 7 \pmod{8} \\ wc_{u^0}(x) + w^2 c_{u^3}(x) + \sum_{i=1,2,4,5} c_{u^i}(x) & \text{if } p = 3 \pmod{8} \end{cases}$$

(c) The trace representation and linear complexity of s is given as follows:

$$s(t) = \sum_{\substack{0 \le m < c \\ m \equiv 0 \pmod{6}}} \operatorname{Tr}_{1}^{n} \left(\beta^{u^{m}t}\right) = \sum_{\substack{m=0 \\ m=0}}^{c/6-1} \operatorname{Tr}_{1}^{n} \left(\beta^{u^{6}mt}\right), \quad LC = (p-1)/6,$$

$$m \equiv 0 \pmod{6}$$

$$s(t) = \sum_{\substack{0 \le m < c \\ m \equiv 0 \pmod{6}}} \operatorname{Tr}_{1}^{n} \left(\omega\beta^{u^{m}t}\right) + \sum_{\substack{0 \le m < c \\ m \equiv 3 \pmod{6}}} \operatorname{Tr}_{1}^{n} \left(\omega^{2}\beta^{u^{m}t}\right) + \sum_{\substack{0 \le m < c \\ m \not\equiv 0 \pmod{6}}} \operatorname{Tr}_{1}^{n} \left(\beta^{u^{m}t}\right), \quad LC = p-1.$$

### **Theorem 5** Let p = ef + 1 with e = 4 and f odd. Then

- 1. There exists a generator u of  $F_p^*$  with  $2 \in uH_4$  and  $\beta \in <\alpha >^*$ , such that  $c_{u^i}(\beta) = (\theta, \theta^2, \theta^4, \theta^8)$ , where  $\theta = \rho$  or  $\rho + 1$ , and  $\rho$  is a root of the polynomial  $x^4 + x^3 + 1$  and is a primitive 15-th root of unity, and hence,  $\rho + 1$  is a root of the polynomial  $\sum_{0 \leq i \leq 4} x^i$  and is a primitive 5-th root of unity.
- 2. In case when  $p = 4f + 1 = 1 + 4z^2$  for some integer z (for this case, it is known that  $H_4$  is a (p, (p-1)/4, (p-5)/16)-cyclic difference set mod p), let s be the characteristic sequence of  $H_4$ . Then  $s = \underline{1} + \mathbf{b}_{u^0}$ , and it has a defining pair  $(g(x), \beta)$ , where

$$g(x) = \sum_{0 \le i < 4} \theta^{2^{2+i}} c_{u^i}(x),$$

and  $\theta$  is described as in the item 1 above. As a consequence,  $LC(\mathbf{s}) = p - 1$ .

3. In case when  $p = 9 + 4z^2$  for some integer z (for this case, it is known that  $H_4 \cup \{0\}$  is a (p, (p+3)/4, (p+3)/16)- cyclic difference set mod p), let s be the characteristic sequence of the difference set  $H_4 \cup \{0\}$ . Then  $s = \underline{1} + \mathbf{b}_* + \mathbf{b}_{u^0}$ , and it has a defining pair  $(g(x), \beta)$ , where

$$g(x) = 1 + \sum_{0 \le i < 4} (\theta^{2^{2+i}} + 1) c_{u^i}(x),$$

and  $\theta$  is described as in the item 1 above. As a consequence,  $LC(\mathbf{s}) = p$ .

**Theorem 6** Let p = ef + 1 with e = 8 and f odd, and assume d = 8, where d is the d-parameter corresponding to (p, e). Then

1. There exist u and  $\beta \in <\alpha >^*$  such that  $\mathbf{c}_u(\beta) = (c_0, c_1, \cdots, c_7)$ , where

 $(c_0, c_1, \cdots, c_7) = (1, 1, 0, 1, 0, 0, 0, 0),$  or its complement (0, 0, 1, 0, 1, 1, 1, 1).

2. In the case when  $p = 1+8z^2 = 9+64y^2$  for some odd integers z and y (for this case, it is known that  $H_8$  is a (p, (p-1)/8, (p-7)/64)-cyclic difference set mod p), let s be the characteristic sequence of  $H_8$ . Then  $s = \underline{1} + \mathbf{b}_{u^0}$ , and it has a defining pair  $(g(x), \beta)$ , where

$$g(x) = \sum_{0 \le i < 8} c_{4+i} c_{u^i}(x),$$

the indexes 4 + i is modulo 8, and  $c_i$  is described as in the item 1 above.

3. In the case when  $p = 49 + 8z^2 = 441 + 64y^2$  for some odd integers z and y (for this case, it is known that  $D = H_8 \cup \{0\}$  is a (p, (p+7)/8, (p+7)/64)-cyclic difference set mod p), let s be the characteristic sequence of  $D = H_8 \cup \{0\}$ . Then  $s = \underline{1} + b_* + b_{u^0}$ , and it has a defining pair  $(g(x), \beta)$ , where

$$g(x) = 1 + \sum_{0 \le i < 8} (c_{4+i} + 1)c_{u^i}(x),$$

the subscript 4 + i is computed mod 8, and  $c_i$  is described as in the item 1 above.

**Theorem 7** Let p = 31, e = 10, and let **s** be the characteristic sequence of the cyclic difference set  $D = H_{10} \cup 11H_{10} = \{i \pmod{31} \mid i = 1, 5, 11, 24, 25, 27\}$ . Let  $\beta$  be a root of the polynomial  $x^5 + x^2 + 1$ . Then 1.  $\mathbf{c}_{11}(\beta) = (c_0, c_1, \dots, c_9)$ , where  $c_{2j} = \beta^{-7 \cdot 2^{4j}}$ ,  $c_{2j+1} = \beta^{-2^{4j}}$ ,  $0 \le j < 5$ . 2.  $\mathbf{s} = \underline{1} + \mathbf{b}_1 + \mathbf{b}_{11}$ .

3. *Let* 

$$g(x) = 1 + \sum_{0 \le j < 5} \left( \beta^{11 \cdot 2^{4j}} c_{11^{2j}}(x) + \beta^{18 \cdot 2^{4j}} c_{11^{2j+1}}(x) \right)$$

Then  $(g(x), \beta)$  is a defining pair of s.

- Binary sequences (of period p) of all the cyclic difference sets D which are some union of cosets of e-th powers in  $F_p^*$  for  $e \le 12$  are studied in terms of
  - their defining pairs,
  - trace representations,
  - linear complexities.
- In particular, linear complexities of all the *e*-th residue sequences are determined whenever  $d = \gcd(e, (p-1)/n) = 1$ , where *n* is the order of 2 mod *p*.
- How to evaluate the  $e\text{-tuple}~(c_{u^0}(\beta),...,c_{u^{e-1}}(\beta))$  for some u and  $\beta$  whenever a prime p=ef+1 is given ?
- Open Problem: Which one among the two values ρ and ρ + 1 the element θ in Theorem 4 or in Theorem 5 takes has not been determined yet, and we do not know whether both values will be taken when p changes; and the same problem for the tuple (c<sub>0</sub>, c<sub>1</sub>, ..., c<sub>7</sub>) in Theorem 6.