

Properties and Crosscorrelation of Decimated Sidelnikov Sequences**

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SUMMARY In this paper, we show that if the d -decimation of a $(q-1)$ -ary Sidelnikov sequence of period $q-1 = p^m - 1$ is the d -multiple of the same Sidelnikov sequence, then d must be a power of a prime p . Also, we calculate the crosscorrelation magnitude between some constant multiples of d - and d' -decimations of a Sidelnikov sequence of period $q-1$ to be upper bounded by $(d+d'-1)\sqrt{q} + 3$.

key words: Sidelnikov sequences, d -decimation sequences, constant-multiple sequences, correlation bound, Weil bound

1. Introduction

For a prime power $q = p^m$ and a positive integer M such that $M|q-1$, Sidelnikov [1] introduced an M -ary sequences (called the Sidelnikov sequences) of period $q-1$, and showed that their non-trivial autocorrelation magnitudes are upper bounded by 4 regardless of M and q .

Binary and non-binary sequences with good autocorrelation properties have important applications in spread spectrum communications and radar engineering [2], [3]. Sequence family with further property on good (low) cross-correlation has also some important applications in multi-user communications such as cellular mobile communications [3], [4]. Non-binary sequence family constructions using the Sidelnikov sequences have been considered by many researchers [5], [6], [8]–[13].

Sampling or decimation is a well-known method for constructing a new sequence from the given sequence [2]–[4]. In this paper, we show that if the d -decimation of a $(q-1)$ -ary Sidelnikov sequence of period $q-1 = p^m - 1$ is the d -multiple of the same Sidelnikov sequence, then d must be a power of a prime p . Also, we calculate the crosscorrelation magnitude between some constant multiples of d - and d' -

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decimations of a Sidelnikov sequence of period $q-1$ to be upper bounded by $(d+d'-1)\sqrt{q} + 3$.

Section 2 introduces some preliminary concepts for the main results: notation, definition of Sidelnikov sequences, correlation, Weil bound on character sums, and decimation. Main results are given in Sect. 3. Section 4 gives a brief concluding remark.

2. Preliminaries

2.1 Notation and Convention

We will use the following notation:

- p : a prime number
- q : a prime power p^m with positive integer m
- $GF(q)$: the finite field with q elements
- d : a decimation factor with $\gcd(q-1, d) = 1$
- M : a divisor of $q-1$ with $M \geq 2$
- $\omega_M = \exp(j\frac{2\pi}{M})$ where $j = \sqrt{-1}$
- β : a fixed primitive element of $GF(q)$
- ψ : the multiplicative character of $GF(q)$ of order M defined by

$$\psi(x) = \exp\left(j\frac{2\pi}{M} \log_{\beta} x\right) = \omega_M^{\log_{\beta} x}$$

We keep $\log(0) = 0$ and $\psi(0) = 1$ in this paper for convenience.

2.2 Sidelnikov Sequences

Definition 1 [1] For any fixed primitive element β of $GF(q)$, let $D_k = \{\beta^{Mi+k} - 1 \mid 0 \leq i < \frac{q-1}{M}\}$. Then M -ary Sidelnikov sequence $\{s(t)\}$ of period $q-1$ is defined as

$$s(t) = \begin{cases} 0, & \text{if } \beta^t = -1 \\ k, & \text{if } \beta^t \in D_k. \end{cases}$$

We will say that $s(t)$ above is defined by the primitive element β .

Equivalently, it can also be expressed, with the new convention that $\log 0 = 0$, as

$$s(t) \equiv \log_{\beta}(\beta^t + 1) \pmod{M}$$

for $0 \leq t \leq q-2$.

2.3 Correlation

A correlation is a measure of distance between two sequences or sequence families. If two sequences are the same, the correlation is called auto-correlation. Otherwise, we call it cross-correlation. The following definition has been well-known [4].

Definition 2 Let $\{a(t)\}$ and $\{b(t)\}$ be M -ary sequences of period L , where $0 \leq t \leq L-1$. A periodic correlation between $\{a(t)\}$ and $\{b(t)\}$ is defined by

$$C_{a,b}(\tau) = \sum_{t=0}^{L-1} \omega_M^{a(t)-b(t+\tau)}$$

for $0 \leq \tau \leq L-1$, where $\omega_M = \exp(j\frac{2\pi}{M})$ and $t+\tau$ is computed modulo L .

For $\tau \equiv 0 \pmod{L}$ when $\{a(t)\}$ and $\{b(t)\}$ are the same sequence, the value above becomes L , which is regarded as trivial. The values at all other cases (for all τ when $\{a(t)\}$ and $\{b(t)\}$ are two different sequences or for $\tau \not\equiv 0 \pmod{L}$ when $\{a(t)\}$ and $\{b(t)\}$ are the same sequence) are called non-trivial.

For a sequence set \mathcal{S} , $C_{\max}(\mathcal{S})$ is the maximum magnitude of all the non-trivial correlations of the pairs of sequences in \mathcal{S} .

2.4 Weil Bound

The Weil bound gives an upper bound on the multiplicative character sums, and has been used to calculate upper bound on the crosscorrelation of various sequences [4]. Yu and Gong [9], [13] introduced a refined version from some of the previously well-known ones, with an additional assumption that $\psi(0) = 1$. We will use a further refined version by Kim [10].

Theorem 1 [10] Let $f_1(x), \dots, f_m(x)$ be distinct monic irreducible polynomials over $GF(q)$ with degrees d_1, \dots, d_m , with e_j the number of distinct roots in $GF(q)$ of $f_j(x)$ ($j = 1, \dots, m$). Let ψ_1, \dots, ψ_m be nontrivial multiplicative characters of $GF(q)$, with $\psi_j(0) = 1$ ($j = 1, \dots, m$). Then, for $a_1, \dots, a_m \in GF(q)^\times$, we have the estimate

$$\left| \sum_{x \in GF(q)} \psi_1(a_1 f_1(x)) \cdots \psi_m(a_m f_m(x)) \right| \leq \left(\sum_{j=1}^m d_j - 1 \right) \sqrt{q} + \sum_{j=1}^m e_j.$$

2.5 Decimation and Constant Multiple

Taking decimation and/or constant multiple is a well-known method for constructing a new sequence from the given sequence [5], [7].

Definition 3 Let $a(t)$ be an M -ary sequence of period L . Then (1) the d -decimation sequence $b(t)$ of $a(t)$ is $b(t) = a(dt)$ for $t = 0, 1, \dots$; (2) the d -multiple sequence $c(t)$ of $a(t)$ is $c(t) = d \cdot a(t)$ for $t = 0, 1, \dots$

When we apply the d -decimation to the periodic sequence of period L , we see easily that the period of decimated sequence becomes $\frac{L}{k}$ where $k = \gcd(d, L)$. Therefore, we have to choose d such that $\gcd(d, L) = 1$ in order to maintain the original period.

When we apply the d -multiple to the M -ary sequence, we see easily that the number of distinct values of the sequence becomes $\frac{M}{k}$ where $k = \gcd(d, M)$. Therefore, we have to choose d such that $\gcd(d, M) = 1$ in order to keep the resultant sequence to have truly M distinct values, i.e., to be an M -ary sequence.

3. Main Result

The Sidelnikov sequence is constructed using a primitive element of a finite field. It is well-known that if β is a primitive element of $GF(q)$, then β^d is also a primitive element of $GF(q)$ whenever $\gcd(d, q-1) = 1$. It is interesting to note that the d -decimation of a Sidelnikov sequence defined by a primitive element β is a d -multiple of another Sidelnikov sequence defined by the primitive element β^d . Following has been essentially proved as Lemma 2 in [7]:

Theorem 2 (Lemma 2 in [7]) Let $q = p^m$ and $\{s(t)\}$ be an M -ary Sidelnikov sequence of period $q-1$. Let d be relatively prime to $q-1$. Then $s(dt) \equiv d \cdot s'(t) \pmod{M}$ where $\{s'(t)\}$ given by $s'(t) \equiv \log_\gamma(y^t + 1) \pmod{M}$ is an another M -ary Sidelnikov sequence of period $q-1$ defined by the primitive element $\gamma = \beta^d$.

Example 1 Table 1 shows some decimation sequences of a 10-ary Sidelnikov sequence of period $q-1 = 10$. $GF(11)$ has 4 primitive elements: $2, 6 \equiv 2^9, 7 \equiv 2^7, 8 \equiv 2^3$. We can show that $9 \cdot s_1(t) \equiv s(9t)$ where $6 \equiv 2^9$. Similarly, $7 \cdot s_2(t) \equiv s(7t)$ and $3 \cdot s_3(t) \equiv s(3t)$.

Corollary 1 Let $q = p^m$ and $\{s(t)\}$ be an M -ary Sidelnikov sequence of period $q-1$. If $d = p^l k$ with $\gcd(p, k) = 1$ and a nonnegative integer l , then $s(p^l kt) = p^l s(kt)$ for all t . In particular, when $k = 1$, this implies $s(p^l t) = p^l s(t)$ for all t .

The above corollary is quite straightforward since β

Table 1 Decimation sequences of a Sidelnikov sequence in Example 1.

t	0	1	2	3	4	5	6	7	8	9
$\beta = 2$	$s(t)$	1	8	4	6	9	0	5	3	2
$\gamma = 6 \equiv 2^9$	$s_1(t)$	9	3	8	7	5	0	1	4	6
$\gamma = 7 \equiv 2^7$	$s_2(t)$	3	9	7	4	6	0	2	1	5
$\gamma = 8 \equiv 2^3$	$s_3(t)$	7	2	5	9	8	0	4	6	3

and β^{p^l} generate the same Sidelnikov sequence, but its converse is not at all trivial. We prove the converse for the case of $M = q - 1$ and $k = 1$. We guess that it is also true for any divisor M of $q - 1$ and $k = 1$, but leave this as an open problem.

Theorem 3 Let $q = p^m$, and $\{s(t)\}$ be a $(q - 1)$ -ary Sidelnikov sequence of period $q - 1$. If there exists a d such that $s(dt) = d \cdot s(t)$ for all t , then $d = p^l$ for some nonnegative integer l .

proof: Since the period of the sequence is $q - 1$, we may assume that $d \leq q - 2$. We assume that $s(dt) = d \cdot s(t)$ for all t , and write, on the contrary, $d = p^l r$, where l is a nonnegative integer, $\gcd(r, p) = 1$ and $r > 1$. Then we will consider the binomial coefficient $\binom{d}{p^l}$ and see if it is divisible by p or not. It will turn out that it is divisible by p in one way, and is not in other way, giving a desired contradiction to $r > 1$.

First, we have $\log_\beta(\beta^t + 1)^d = \log_\beta(\beta^{dt} + 1)$, and hence $(\beta^t + 1)^d = \beta^{dt} + 1$ for all $0 \leq t \leq q - 2$. Therefore, $(a + 1)^d = a^d + 1$ for all $a \in GF(q)$. This implies that $p \mid \binom{d}{i}$ for all $1 \leq i \leq d - 1$. In particular, this implies that $\binom{d}{p^l}$ is divisible by p .

Now we use Kummer's criterion [14]: the power of p dividing $\binom{n}{k}$ is the number of carries when we add k to $n - k$ in base p . We may divide r by p and write $r = pa + j$ for some $1 \leq j \leq p - 1$. By Kummer's criterion, the power of p dividing $\binom{d}{p^l}$ is the number of carries when adding p^l to $d - p^l = p^l(r - 1)$ in base p . Observe that the right-most non-zero digit (in base p) of $p^l(r - 1)$ is $(j - 1)$ when $a = 0$ and those of p^l is 1, and hence, there are no carries in the sum that is just j . Therefore, $\binom{d}{p^l}$ is not divisible by p . Similarly, it is obvious that $\binom{d}{p^l}$ is not divisible by p when $a > 0$. This is a desired contradiction to $r > 1$. ■

Let $s(t)$ be an M -ary Sidelnikov sequence of period $q - 1$. We will now derive the maximum correlation bound between $c_1 \cdot s(dt)$ and $c_2 \cdot s(d't)$ where c_1, c_2, d and d' are some constants. If p divides d , i.e. $d = p^l k$, where $l > 0$, then $s(dt) = p^l \cdot s(kt)$ by Corollary 1. Especially, if $d = p^l$ and $d' = p^{l'}$, then the correlation between $c_1 \cdot s(dt)$ and $c_2 \cdot s(d't)$ becomes just a correlation between two different constant multiples of a Sidelnikov sequence, and it was computed earlier by Kim and Song [5] and also by Yu and Gong [13]. Consequently, we need to consider only the case where p divides neither d nor d' .

Lemma 1 Let d, d' be positive integers with $(d, q - 1) = (d', q - 1) = 1$, and let $0 \leq \tau \leq q - 2$. Then we have the following:

- (a) The only root in $GF(q)$ of $x^d + 1 = 0$ is -1 . The only root in $GF(q)$ of $x^{d'} + \beta^{-d'\tau} = 0$ is $-\beta^{-\tau}$.
- (b) If p does not divide d , then $x^d + 1$ has no multiple roots. If p does not divide d' , then $x^{d'} + \beta^{-d'\tau}$ has no multiple roots.

(c) $x^d + 1$ and $x^{d'} + \beta^{-d'\tau}$ have a common root if and only if $\tau = 0$.

proof: (a) If γ and δ are roots of $x^d + 1$, then $(\frac{\gamma}{\delta})^d = 1$. But the only root of $x^d = 1$ in $GF(q)$ is 1, as $\beta^{id} = 1 \Rightarrow q - 1 \mid id \Rightarrow q - 1 \mid i$, since $(d, q - 1) = 1$. So $\beta^i = 1$, and hence $\gamma = \delta$. If d is odd, then the root of $x^d + 1 = 0$ is -1 . And if d is even, then q is even since $(d, q - 1) = 1$, and hence the characteristic is 2. Hence, $-1 = +1$ is the only root of $x^d + 1$ over $GF(q)$, q even. The other case is similar.

(b) If p does not divide d , the $x^d + 1$ and its derivative dx^{d-1} are relatively prime. The other case is similar.

(c) If $\tau = 0$, then -1 is the common root. Conversely, let $\gamma^d + 1 = 0$ and $\gamma^{d'} + \beta^{-d'\tau} = 0$. Then $\gamma = -\zeta = -\beta^{-\tau}\eta$, with $\zeta^d = 1$ and $\eta^{d'} = 1$. So $\beta^\tau = \eta\zeta^{-1}$. Raising dd' -th power of both sides, we have $\beta^{\tau dd'} = 1$, which implies $\tau = 0$, since $(d, q - 1) = (d', q - 1) = 1$. ■

Theorem 4 Assume that $(d, q - 1) = (d', q - 1) = 1$ and that p divides neither d nor d' . Let $a(t) = c_1 \cdot s(dt)$ and $b(t) = c_2 \cdot s(d't)$ are cyclically inequivalent for some M -ary Sidelnikov sequence $s(t)$ of period $q - 1$ and constants $1 \leq c_1, c_2 \leq M - 1$. Then we have

$$\left| \max_{\tau} \{C_{a,b}(\tau)\} \right| \leq (d + d' - 1) \sqrt{q} + 3$$

where τ runs over the integers $0 \leq \tau \leq q - 2$.

proof: By the definition of Sidelnikov sequences and its decimation, we see that $a(t) = c_1 \log_\beta(\beta^{dt} + 1)$ and $b(t) = c_2 \log_\beta(\beta^{d't} + 1)$. Then, their correlation becomes as follows:

$$\begin{aligned} C_{a,b}(\tau) &= \sum_{t=0}^{q-2} \omega_M^{a(t)-b(t+\tau)} \\ &= \sum_{t=0}^{q-2} \omega_M^{c_1 \log_\beta(\beta^{dt} + 1) - c_2 \log_\beta(\beta^{d'(t+\tau)} + 1)} \\ &= \sum_{x \in GF(q)} \psi^{c_1}(x^d + 1) \cdot \psi^{M-c_2}(\beta^{d'\tau} x^{d'} + 1) - 1 \\ &= \sum_{x \in GF(q)} \psi^{c_1}(x^d + 1) \cdot \psi^{M-c_2}(\beta^{d'\tau} (x^{d'} + \beta^{-d'\tau})) - 1 \end{aligned}$$

Case 1. $\tau \neq 0$.

Lemma 1 says that, for any τ with $1 \leq \tau \leq q - 2$, $f_1(x) = x^d + 1$ and $f_2(x) = x^{d'} + \beta^{-d'\tau}$ are relatively prime (cf. Lemma 1(c)) with the respective number e, e' of distinct roots in $GF(q)$ equal to $e = e' = 1$. (cf. Lemma 1(a)). Since M does not divide c_1 and $M - c_2$, $\psi_1 = \psi^{c_1}$ and $\psi_2 = \psi^{M-c_2}$ are not trivial. Therefore we have

$$\begin{aligned} &\psi_1(x^d + 1) \cdot \psi_2(\beta^{d'\tau} (x^{d'} + \beta^{-d'\tau})) \\ &= \psi_1(x+1) \cdot \psi_1(h_1(x)) \cdot \psi_2(\beta^{d'\tau} (x+\beta^{-\tau})) \cdot \psi_2(h_2(x)) \end{aligned}$$

where $x^d + 1 = (x+1)h_1(x)$ and $x^{d'} + \beta^{-d'\tau} = (x+\beta^{-\tau})h_2(x)$ for some polynomials $h_1(x)$ and $h_2(x)$. By Lemma 1, $h_1(x)$ and

Table 2 Comparison between true max and bound in Cor. 2 for $c = 1$, $d' = 1$, $\gcd(d, q - 1) = 1$ and p not dividing d .

p	q	d	M	True Max	Bound = $d\sqrt{q} + 3$
2	$64 = 2^6$	5	7	17.62	43.00
3	$243 = 3^5$	5	11	41.78	80.94
2	$256 = 2^8$	7	15	49.79	115.00
17	$289 = 17^2$	5	8	45.96	88.00
		7	8	38.88	122.00
7	$343 = 7^3$	5	9	47.78	95.60
2	$512 = 2^9$	5	7	56.58	68.88

$h_2(x)$ are products of some distinct monic irreducible polynomials over $GF(q)$ of degrees greater than 1 (cf. Lemma 1(b)). Therefore, we can apply Weil bound in Theorem 1, and hence

$$|C_{a,b}(\tau)| \leq (d + d' - 1)\sqrt{q} + 3.$$

Case 2. $\tau = 0$.

In this case, we have

$$C_{a,b}(\tau = 0) = \sum \psi^{c_1}(x^d + 1) \cdot \psi^{M-c_2}(x^{d'} + 1) - 1. \quad (1)$$

Assume that $d = d'$. Then (1) becomes $\sum \psi^{c_1-c_2}(x^d + 1) - 1$. If $c_1 = c_2$, then two sequences are the same. Otherwise,

$$|C_{a,b}(\tau = 0)| \leq (d - 1)\sqrt{q} + 2$$

by Weil bound and Lemma 1.

Otherwise, assume that $d \neq d'$. Then (1) becomes $\sum \psi^{c_1-c_2}(x+1) \cdot \psi^{c_1}(x^{d-1} - x^{d-2} + \dots - x + 1) \cdot \psi^{M-c_2}(x^{d'-1} - x^{d'-2} + \dots - x + 1) - 1$. Second and third polynomials of degree $d-1$ and $d'-1$ inside the character are products of distinct monic irreducible polynomials over $GF(q)$ by Lemma 1. Therefore, we can apply Weil bound, and hence

$$|C_{a,b}(\tau = 0)| \leq (d + d' - 2)\sqrt{q} + 2.$$

Corollary 2 Assume that $\gcd(d, q - 1) = 1$, and that p does not divide d . Let $s(t)$ be an M -ary Sidelnikov sequence of period $q - 1$. Let $a(t) = c \cdot s(t)$ and $b(t) = s(dt)$. Then we have

$$\left| \max_{\tau} \{C_{a,b}(\tau)\} \right| \leq d\sqrt{q} + 3.$$

Example 2 Table 2 shows the difference between the exact maximal correlation magnitude and our correlation bound given in the above corollary for some $q, d, c = 1$ and M .

Example 3 Figure 1 shows the correlation function of the 11-ary Sidelnikov sequence of period $3^5 - 1 = 242$ and its 5-decimation. The horizontal dotted line in the figure indicates the correlation bound in the main result, also shown

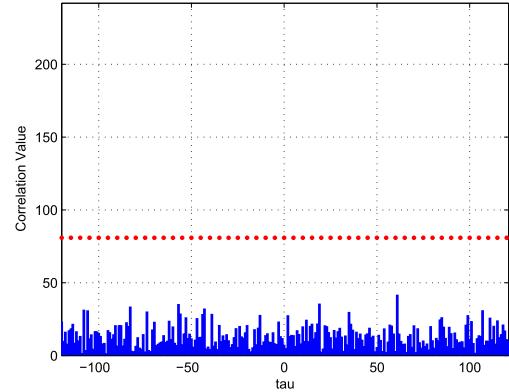


Fig. 1 Correlation between a 5-ary Sidelnikov sequence of period 242 and its 5-decimation.

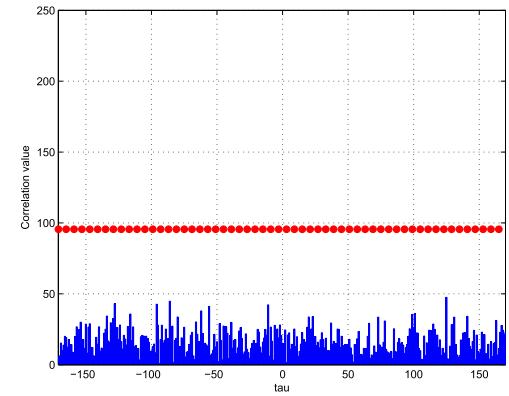


Fig. 2 Correlation between a 9-ary Sidelnikov sequence of period 342 and its 5-decimation.

in Table 2, which is about 81. The true max turns out to be about 42, showing some gap between the two numbers. Figure 2 shows the correlation function of the 9-ary Sidelnikov sequence of period $7^3 - 1 = 342$ and its 5-decimation. In this case, the true max turns out to be about 48.

4. Concluding Remark

Both figures in the above example show that there is a gap between the true max and the bound calculated using the Weil bound. This gap will increase when the constant multiple c is not relatively prime to M . For example, when $q = 343$ and $M = 9$ (cf. Example for Fig. 2), the true max between $3 \cdot s(t)$ and $s(5t)$ is around 20 while the bound in Cor. 2 is also 95.6. We guess that these gaps can be reduced by some direct calculation of the correlation between $c \cdot s(t)$ and $s(dt)$ for $1 \leq c \leq M - 1$ and $\gcd(d, q - 1) = 1$. It would be interesting in the future to see and compare these gaps between those from the direct calculation and from the Weil bound.

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