



Chaotic Behavior of Integer Sequences from Linear Feedback Shift Registers

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➤ Introduction

- What is chaotic behavior?
- Problems in Discrete Chaos
- Discrete Lyapunov Exponent of Discrete Chaos

➤ Main Contribution

- Chaotic behavior of integer sequences from LFSRs
- Main Theorem and Proof

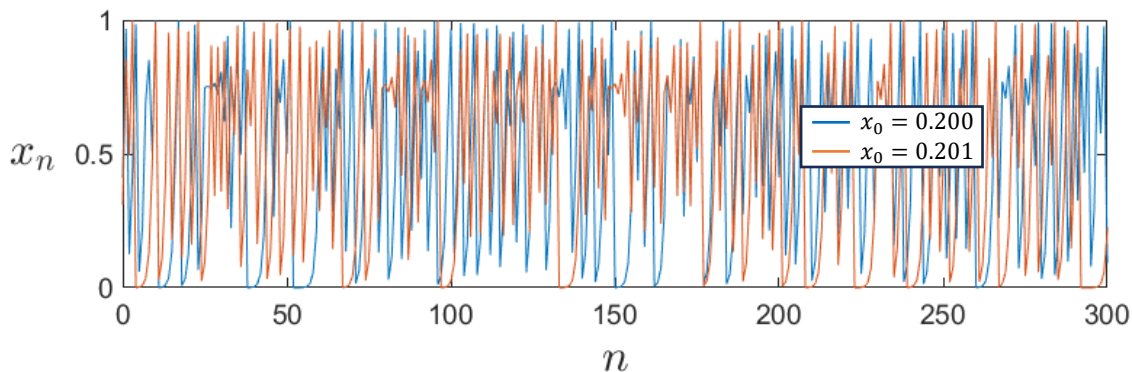
➤ Conclusion



What is chaotic behavior?

- **Chaotic behavior** : aperiodic and highly unpredictable dynamics
- **Chaotic maps** are nonlinear functions that are **extremely sensitive** to initial values, where even small differences can produce **completely different** sequences.

e.g., Logistic maps



- **Applications:**

CRYPTOGRAPHY,
DIGITAL WATERMARKING,
SPREAD-SPECTRUM COMMUNICATION,
RANDOM NUMBER GENERATION,.....



Examples ?

from WIKIPEDIA



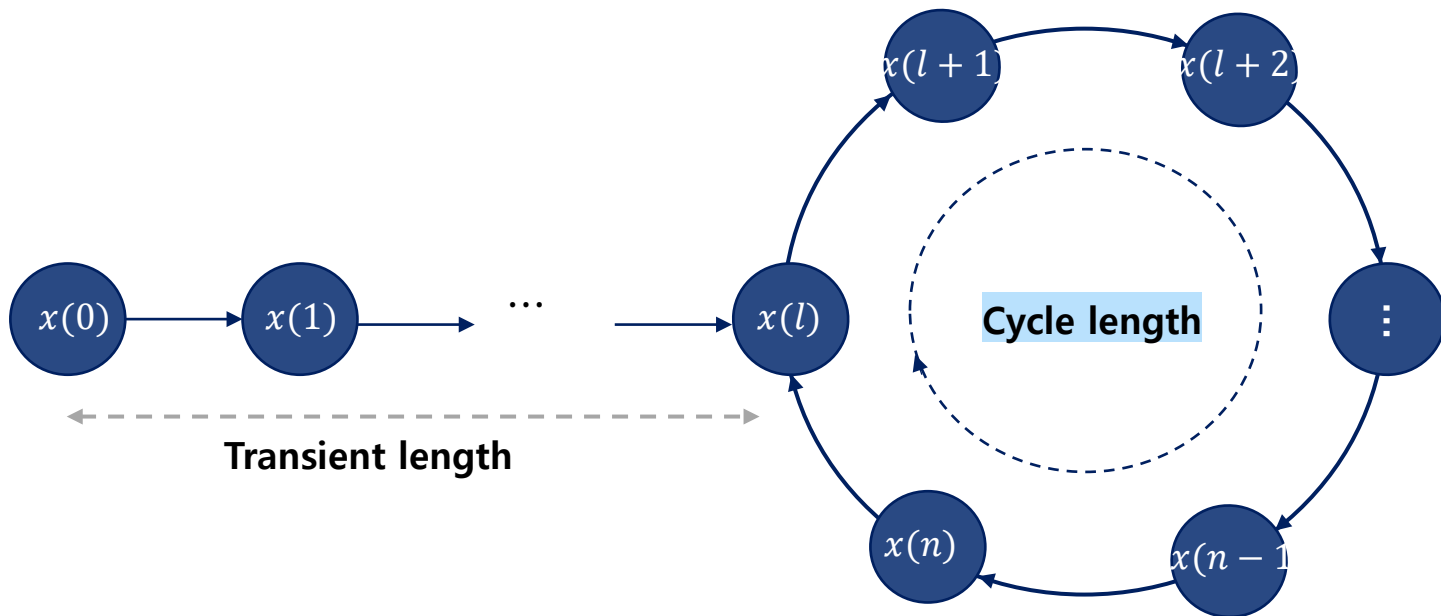
List of chaotic maps				
Map	Time domain	Space domain	Number of space dimensions	Number of parameters
2D Lorenz system ^[1]	discrete	real	2	1
2D Rational chaotic map ^[2]	discrete	rational	2	2
3-cells CNN system	continuous	real	3	
ACT chaotic attractor ^[3]	continuous	real	3	
Aizawa chaotic attractor ^[4]	continuous	real	3	5
Arneodo chaotic system ^[5]	continuous	real	3	
Arnold's cat map	discrete	real	2	0
Baker's map	discrete	real	2	0
Basin chaotic map ^[6]	discrete	real	2	1
Beta Chaotic Map ^[7]				12
Bogdanov map	discrete	real	2	3
Brusselator	continuous	real	3	
Burke-Shaw chaotic attractor ^[8]	continuous	real	3	2
Chen chaotic attractor ^[9]	continuous	real	3	3
Chen-Celikovsky system ^[10]	continuous	real	3	
Chen-LU system ^[11]	continuous	real	3	3
Chen-Lee system	continuous	real	3	
Chialvo map	discrete	discrete		3
Chossat-Golubitsky symmetry map				
Chua circuit ^[12]	continuous	real	3	3
Circle map	discrete	real	1	2
Complex quadratic map	discrete	complex	1	1
Complex squaring map	discrete	complex	1	0
Complex cubic map	discrete	complex	1	2
Clifford fractal map ^[13]	discrete	real	2	4
Degenerate Double Rotor map				
De Jong fractal map ^[14]	discrete	real	2	4
Delayed-Logistic system ^[15]	discrete	real	2	1
Discretized circular Van der Pol system ^[16]	discrete	real	2	1
Discretized Van der Pol system ^[17]	discrete	real	2	2
Double rotor map				
Duffing map	discrete	real	2	2

Problems in Discrete Chaos

When **implemented** in digital systems,
any (real field) chaotic map will suffer from some serious
dynamic degradation

due to **finite-precision** effects:

shorter periods, convergence to fixed values, and
other problems from rounding-off and quantization



Typical orbit of discrete chaotic maps

Discrete Lyapunov Exponent of Discrete chaos

- **Discrete Lyapunov Exponent (DLE) [10,11]**
quantifies the **average rate of divergence** between *neighboring elements* of a permutation F in the discrete domain

$$S := \{1, 2, 3, \dots, M-1, M\}$$



Discrete Lyapunov Exponent of Discrete chaos

➤ Discrete Lyapunov Exponent (DLE) [10,11]

A **permutation** F on a discrete domain $\{1,2,\dots,M\}$
is said to be **discretely chaotic**

if its DLE λ_F satisfies

$$\lim_{M \rightarrow \infty} \lambda_F > 0$$

[10] L. Kocarev and J. Szczepanski, "Finite-space Lyapunov exponents and pseudochaos," *Phys. Rev. Lett.*, vol. 93, p. 234101, **2004**.

[11] L. Kocarev, J. Szczepanski, J.M. Amigo, and I. Tomovski, "Discrete chaos-I: Theory," *IEEE Trans. Circuits Syst. I, Reg. Papers*, vol. 53, no.6, pp.1300–1309, **2006**.



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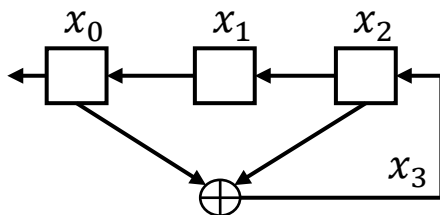
➤ Main Contribution

- Chaotic behavior of integer sequences from LFSRs
- Main Theorem and Proof

➤ Conclusion



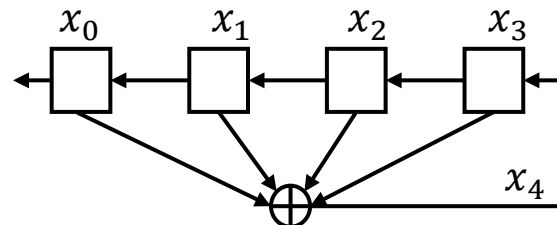
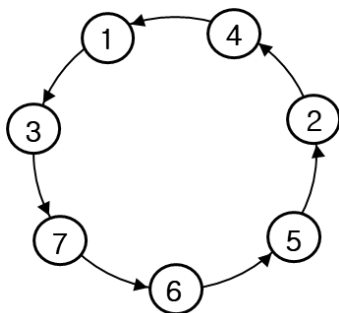
Permutation from the integer sequences of LFSRs



Primitive Connection

Polynomial :

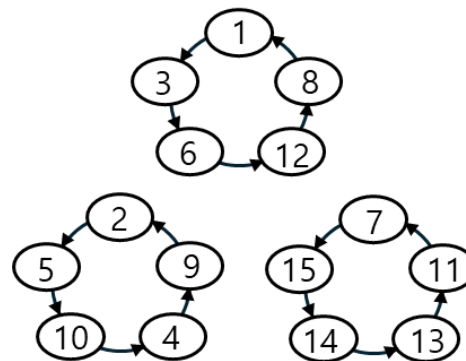
$$x^3 + x^2 + 1$$



Non-primitive but irreducible

Connection Polynomial :

$$x^4 + x^3 + x^2 + x + 1$$



The resulting integer sequences
partition

the set $\{1, 2, \dots, 2^L - 1\}$ into **disjoint cycles** and
induce

a **permutation on the set** $\{1, 2, \dots, 2^L - 1\}$

Is this a **discrete Chaos** ?
Can we calculate its **dLE** ?

Main Theorem

Theorem 1

The discrete Lyapunov Exponent λ_F of F satisfies the following:

$$\ln(\sqrt{3}) \leq \lim_{L \rightarrow \infty} \lambda_F \leq \ln(2),$$

where F is the induced integer permutation from an L -stage LFSR with **irreducible connection polynomial** and **a non-zero initial state**.



*Satisfies the definition of **discrete chaos** !*

DLE λ_F of a permutation F on the set S :

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln(F(z_i) - F(z_{i+1}))$$

where $z_i = i$ for $i \in S$ and $z_{M+1} = M - 1$ [11].

$$S := \{1, 2, 3, \dots, M-1, M\}$$

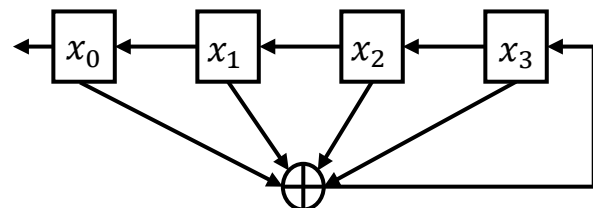
neighbors

$$S := \{1, 2, \dots, 2^L - 1\}$$

$$F: S \rightarrow S$$

z	$F(z)$
$x_0 x_1 x_2 x_3$	$x_0 x_1 x_2 x_3$
1 = 0 0 0 1	0 0 1 1
2 = 0 0 1 0	0 1 0 1
3 = 0 0 1 1	0 1 1 0
0 1 0 0	1 0 0 1
$\vdots$	$\vdots$
0 1 0 1	1 0 1 0
0 1 1 0	1 1 0 0
$2^{L-1} - 1 = 0 1 1 1$	1 1 1 1
$2^{L-1} = 1 0 0 0$	0 0 0 1
1 0 0 1	0 0 1 0
$\vdots$	$\vdots$
1 0 1 0	0 1 0 0
1 0 1 1	0 1 1 1
1 1 0 0	1 0 0 0
1 1 0 1	1 0 1 1
1 1 1 0	1 1 0 1
$2^L - 1 = 1 1 1 1$	1 1 1 0

We will use this running example and prove the result in general



irreducible $x^4 + x^3 + x^2 + x + 1$

We divide $S := \{1, 2, \dots, 2^L - 1\}$ into **three cases**:

Case 1) $z \in D := S \setminus \{2^{L-1} - 1, 2^{L-1}, 2^L - 1\}$

Case 2) $z = 2^{L-1} - 1$

Case 3) $z \in \{2^{L-1}, 2^L - 1\}$



Sketch of Proof

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln |\mathbf{F}(\mathbf{z}_{i+1}) - \mathbf{F}(\mathbf{z}_i)|$$

$$S := \{1, 2, \dots, 2^L - 1\}$$

$$F: S \rightarrow S$$

	z	$F(z)$
	$x_0 x_1 x_2 x_3$	$x_0 x_1 x_2 x_3$
1 =	0 0 0 1	
2 =	0 0 1 0	
3 =	0 0 1 1	
	0 1 0 0	
$\vdots$	0 1 0 1	
	0 1 1 0	
$2^{L-1} - 1 =$	0 1 1 1	
$2^{L-1} =$	1 0 0 0	
	1 0 0 1	
$\vdots$	1 0 1 0	
	1 0 1 1	
	1 1 0 0	
	1 1 0 1	
	1 1 1 0	
$2^L - 1 =$	1 1 1 1	

$2z$

or

$2z + 1$

Case 1) $D := S \setminus \{2^{L-1} - 1, 2^{L-1}, 2^L - 1\}$

Claim 1-(1) For all $z \in D$

$$|F(z+1) - F(z)| \in \{1, 2, 3\}.$$

Update rule:

when the feedback is 0

$$F(z) = 2z$$

$$F(z+1) = 2(z+1)$$

when the feedback is 1

$$F(z) = 2z + 1$$

$$F(z+1) = 2(z+1) + 1$$

Therefore,

$$|F(z+1) - F(z)| = \begin{cases} |2(z+1) - 2z| = 2 \\ |(2(z+1) + 1) - 2z| = 3 \\ |2(z+1) - (2z + 1)| = 1 \\ |(2(z+1) + 1) - (2z + 1)| = 2 \end{cases}$$

Sketch of Proof

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln |\mathbf{F}(\mathbf{z}_{i+1}) - \mathbf{F}(\mathbf{z}_i)|$$

$$S := \{1, 2, \dots, 2^L - 1\}$$

$$F: S \rightarrow S$$

z		$F(z)$	
$x_0 x_1 x_2 x_3$		$x_0 x_1 x_2 x_3$	
1 =	0 0 0 1	0 0 1 1	
2 =	0 0 1 0	0 1 0 1	2
3 =	0 0 1 1	0 1 1 0	1
always not	0 1 0 0	1 0 0 1	3
$\vdots$	0 1 0 1	1 0 1 0	1
	0 1 1 0	1 1 0 0	2
$2^{L-1} - 1 =$	0 1 1 1	1 1 1 1	
$2^{L-1} =$	1 0 0 0	0 0 0 1	
	1 0 0 1	0 0 1 0	
$\vdots$	1 0 1 0	0 1 0 0	2
always	1 0 1 1	0 1 1 1	3
contribute	1 1 0 0	1 0 0 0	1
to the	1 1 0 1	1 0 1 1	3
feedback	1 1 1 0	1 1 0 1	2
$2^L - 1 =$	1 1 1 1	1 1 1 0	

Case 1) $D := S \setminus \{2^{L-1} - 1, 2^{L-1}, 2^L - 1\}$

Claim 1-(2) For all $z \in D$

$$|\{z \in D \mid |F(z+1) - F(z)| = 3\}|$$

$$= |\{z \in D \mid |F(z+1) - F(z)| = 1\}|$$

For $z \in [1, 2^{L-1} - 1]$,

- $|F(z+1) - F(z)| = 3 \Leftrightarrow |F(z + 2^{L-1} + 1) - F(z + 2^{L-1})| = 1$
- $|F(z+1) - F(z)| = 1 \Leftrightarrow |F(z + 2^{L-1} + 1) - F(z + 2^{L-1})| = 3$



Sketch of Proof

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln |F(z_{i+1}) - F(z_i)|$$

$$S := \{1, 2, \dots, 2^L - 1\}$$

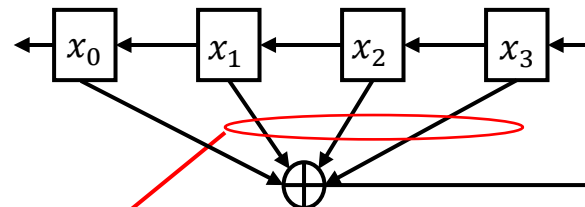
$$F: S \rightarrow S$$

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1 = 0 0 0 1		0 0 1 1
2 = 0 0 1 0		0 1 0 1
3 = 0 0 1 1		0 1 1 0
⋮		⋮
$2^{L-1} - 1 = 0 1 1 1$	→	1 1 1 1 = $2^L - 1$
$2^{L-1} = 1 0 0 0$	→	0 0 0 1 = 1
⋮		⋮
$2^L - 1 = 1 1 1 1$		1 1 1 0

Case 2) For $z = 2^{L-1} - 1$,

Claim 2: $|F(2^{L-1}) - F(2^{L-1} - 1)| = 2^L - 2$

An irreducible polynomial has an **odd number of terms** including x^L and x^0 .



odd number of terms
(except for the first and the last)
contribute to the feedback
FOR ANY IRRED. POLYNOMIAL



Sketch of Proof

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln |\mathbf{F}(\mathbf{z}_{i+1}) - \mathbf{F}(\mathbf{z}_i)|$$

$$\mathcal{S} := \{1, 2, \dots, 2^L - 1\}$$

$$F: \mathcal{S} \rightarrow \mathcal{S}$$

z	$F(z)$
$x_0 x_1 x_2 x_3$	$x_0 x_1 x_2 x_3$
1 = 0 0 0 1	
2 = 0 0 1 0	
3 = 0 0 1 1	
0 1 0 0	
$\vdots$ 0 1 0 1	
0 1 1 0	
$2^{L-1} - 1 = 0 1 1 1$	
$2^{L-1} = 1 0 0 0$	$\rightarrow 0 0 0 1 = 1$
1 0 0 1	$\rightarrow 0 0 1 x \rightarrow$
$\vdots$ 1 0 1 0	
1 0 1 1	
1 1 0 0	
1 1 0 1	
1 1 1 0	
$2^L - 1 = 1 1 1 1$	1 1 0 x $\leftarrow$ similarly
	1 1 1 0

Case 3) for $z = 2^{L-1}$ or $z = 2^L - 1$.

Claim 3: Depending on whether the term x^{L-1} appears **or not** in the polynomial, we have,

$$|F(z + 1) - F(z)| \in \{1, 2\}.$$

when x^{L-1} appears in the conn. poly.

$$0 0 1 x = 0 0 1 0 = 2$$

$$\Rightarrow |F(z + 1) - F(z)| = 1$$

when x^{L-1} DOES NOT appear in the conn. poly.

$$0 0 1 x = 0 0 1 1 = 3$$

$$\Rightarrow |F(z + 1) - F(z)| = 2$$



Sketch of Proof

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln |F(z_{i+1}) - F(z_i)|$$

diff

As a summary,

when x^{L-1} appears in the conn. poly.

$$|\{z \in S \mid \text{diff} = 2\}| = \beta \rightarrow \beta \ln 2$$

$$|\{z \in S \mid \text{diff} = 3\}| = \alpha \rightarrow \alpha \ln 3$$

$$|\{z \in S \mid \text{diff} = 1\}| = \alpha + 2 \rightarrow 0$$

$$|\{z \in S \mid \text{diff} = 2^L - 2\}| = 1 \rightarrow 1 \ln(2^L - 2)$$

$$(2^L - 1) \lambda_F$$

z	$F(z)$
$x_0 x_1 x_2 x_3$	$x_0 x_1 x_2 x_3$
1 = 0 0 0 1	0 0 1 1
2 = 0 0 1 0	0 1 0 1
3 = 0 0 1 1	0 1 1 0
	0 1 0 0
$\vdots$	0 1 0 1
	0 1 1 0
$2^{L-1} - 1 = 0 1 1 1$	1 1 1 1
$2^{L-1} = 1 0 0 0$	0 0 0 1
	1 0 0 1
$\vdots$	1 0 1 0
	1 0 1 1
	1 1 0 0
	1 1 0 1
	1 1 1 0
$2^L - 1 = 1 1 1 1$	1 1 1 0

(claim 1)
1 or 2 or 3

(claim 2) $2^L - 2$

(claim 3) 1 or 2

when x^{L-1} **DOES NOT** appear in the conn. poly.
similar calculation follows.



Sketch of Proof

$$\lambda_F = \frac{1}{M} \sum_{i=1}^M \ln |\mathbf{F}(\mathbf{z}_{i+1}) - \mathbf{F}(\mathbf{z}_i)|$$

$$\lambda_F = \frac{\alpha \ln 3 + \beta \ln 2 + \ln(2^L - 2)}{2^L - 1} = \frac{\ln(3^\alpha 2^\beta) + \ln(2^L - 2)}{2^L - 1}$$

We use $\sqrt{3} < 2$ for **lower bound** and $3 < 2^2$ for **upper bound**:

$$\frac{\ln(3^\alpha 3^{\beta/2}) + \ln(2^L - 2)}{2^L - 1} < \frac{\ln(3^\alpha 2^\beta) + \ln(2^L - 2)}{2^L - 1} < \frac{\ln(2^{2\alpha} 2^\beta) + \ln(2^L - 2)}{2^L - 1}$$

Or

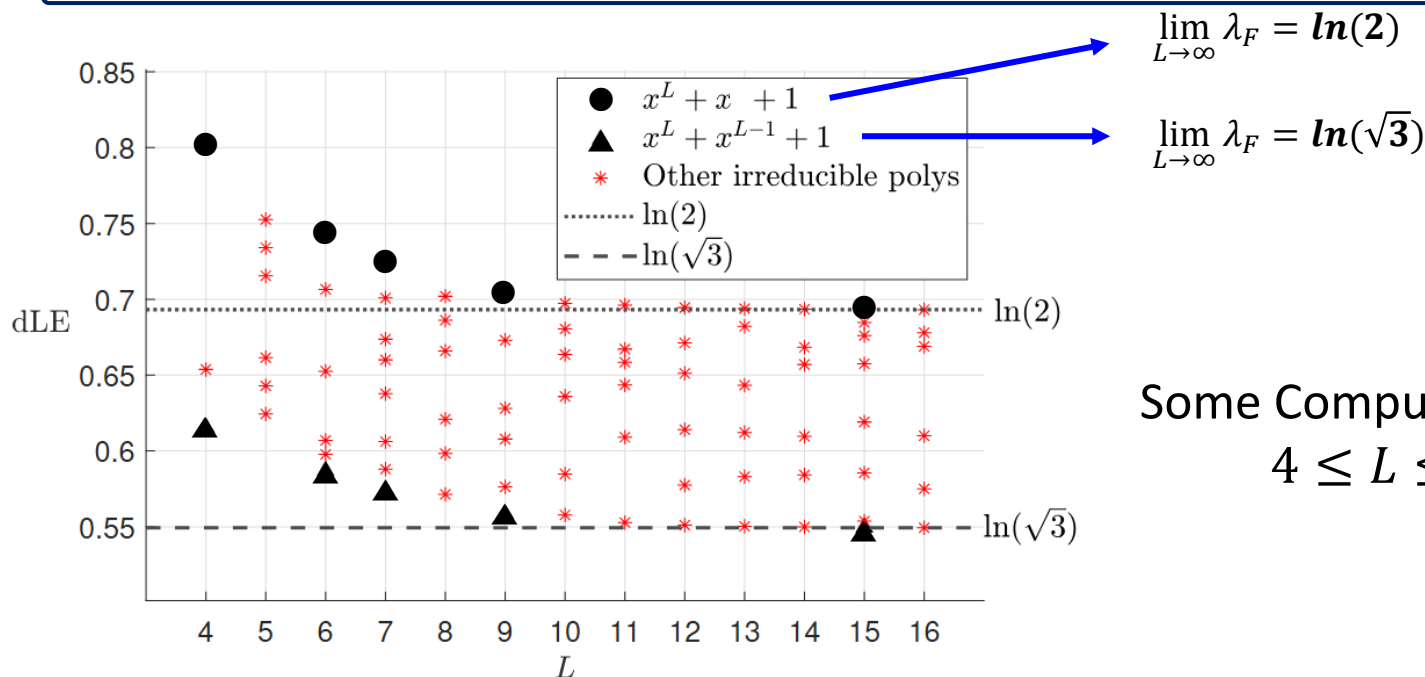
$$\frac{(2\alpha + \beta) \ln(\sqrt{3}) + \ln(2^L - 2)}{2^L - 1} < \lambda_F < \frac{(2\alpha + \beta) \ln(2) + \ln(2^L - 2)}{2^L - 1}$$

We use finally $2\alpha + \beta + 3 = 2^L - 1$ and $L \rightarrow \infty$

$$\ln(\sqrt{3}) \leq \lim_{L \rightarrow \infty} \lambda_F \leq \ln(2)$$

Theorem 1.

$$\ln(\sqrt{3}) \leq \lim_{L \rightarrow \infty} \lambda_F \leq \ln(2)$$



Some Computations for
 $4 \leq L \leq 16$

Conclusion

- Studied the **chaotic nature** of the integer sequences from L -stage **irreducible** LFSRs
- Derived the **discrete Lyapunov Exponent** of **the integer permutation** and showed that it is **lower bound by $\sqrt{3}$** as the register length L **increases indefinitely**.